## Model Checking - Exercise sheet 3

A long time ago, in a galaxy far far away... this exercise sheet was found. Solve the questions and bring balance to the force.

## Exercise 3.4

Let $\varphi=\mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ and $\psi=\neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ be LTL formulas over the atomic propositions $A P=\{p, q, r, s\}$. Say whether the following sequences satisfy $\varphi$ and $\psi$. Justify your answers.
(a) $\emptyset^{\omega}$
(d) $\{r\} \emptyset\{p, q, s\}^{\omega}$
(b) $\{p, q, r, s\}^{\omega}$
(e) $\{r\} \emptyset\{p\}\{q, r\}(\{p, s\} \emptyset)^{\omega}$
(c) $\{p\}^{\omega}$
(f) $\{q, r\} \emptyset\{p, q\} \emptyset\{r, s\}^{\omega}$

## Exercise 3.5

Let $A P=\{s, r, g\}$ be actions of a process: sending a message, receiving a message, and giving a result, respectively. Specify the following properties in LTL, and give example sequences that satisfy and violate the formulas.
(a) The process always gives a result.
(b) The process stops communicating after giving its result.
(c) The process only gives a result once.
(d) The process does nothing until it receives a message.

## Exercise 3.6

Let $A P=\{p, q\}$. An LTL formula is a tautology if it is satisfied by all sequences over $2^{A P}$. Which of the following LTL formulas are tautologies? Justify each answer with a counterexample or a proof.
(a) $\mathbf{G} p \rightarrow \mathbf{F} p$
(d) $\neg \mathbf{F} p \rightarrow \mathbf{F} \neg \mathbf{F} p$
(b) $\mathbf{G}(p \rightarrow q) \rightarrow(\mathbf{G} p \rightarrow \mathbf{G} q)$
(e) $\neg(p \mathbf{U} q) \leftrightarrow(\neg p \mathbf{U} \neg q)$
(c) $\mathbf{F G} p \vee \mathbf{F G} \neg p$
(f) $(\mathbf{G} p \rightarrow \mathbf{F} q) \leftrightarrow(p \mathbf{U}(p \vee q))$

## Exercise 3.1

There are two traffic lights at a road intersection, each of them can be in the following states: $\{$ red, green $\}$. Does the formula $\mathbf{G}\left[\left(t_{1}=r e d \wedge t_{2}=\right.\right.$ green $) \vee\left(t_{1}=\right.$ green $\left.\left.\wedge t_{2}=r e d\right)\right]$ specifies that 'both of the lights should not be green at a given time'? If it does, give an accepting run, otherwise give a counter-example and the correct formula.

## Exercise 3.2

For the model in the last question, write an LTL formula that says 'Both the lights becomes green infinitely many times'.

## Exercise 3.3

You are on your quest to bring balance to the force, for this you have to do some tasks. Look at the following atomic propositions:

- $f$ : Found a death star.
- $d$ : Destroy a death star.
- $k$ : Kill Darth Vader.
- $a$ : You are alive.

Write an LTL formula which specifies that you save the galaxy i.e. killing Darth Vader before dying and whenever you find a death star, destroy it.

## Solution 3.4

(a) - $\emptyset^{\omega} \vDash \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $\emptyset^{\omega} \not \vDash \mathbf{F G} p$ which follows from the fact that $p$ does not occur at all.

- $\emptyset^{\omega} \not \vDash \neg(r \mathbf{U ~ X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $q$ never holds.
(b) - $\{p, q, r, s\}^{\omega} \vDash \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $p$ always occurs and $q$ occurs infinitely often.
- $\{p, q, r, s\}^{\omega} \nLeftarrow \neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $\neg \mathbf{X X} s$ never holds.
(c) • $\{p\}^{\omega} \not \vDash \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $\{p\}^{\omega} \vDash \mathbf{F G} p$ but $\{p\}^{\omega} \not \vDash \mathbf{G F}(q \vee r)$. The former follows from the fact that $p$ occurs infinitely often, and the latter from the fact that $q$ and $r$ never occur.
- $\{p\}^{\omega} \notin \neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $q$ never occurs.
(d) - $\{r\} \emptyset\{p, q, s\}^{\omega} \vDash \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $p$ eventually always occurs and $q$ occurs infinitely often.
- $\{r\} \emptyset\{p, q, s\}^{\omega} \not \vDash \neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $(q \wedge \neg \mathbf{X X} s)$ never holds.
(e) • $\{r\} \emptyset\{p\}\{q, r\}(\{p, s\} \emptyset)^{\omega} \models \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $p$ does not occur eventually always.
- $\{r\} \emptyset\{p\}\{q, r\}(\{p, s\} \emptyset)^{\omega} \not \models \neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $\neg(r \mathbf{U} \mathbf{X} p)$ and $q \wedge \neg \mathbf{X X} s$ does not hold at the first position.
(f) • $\{q, r\} \emptyset\{p, q\} \emptyset\{r, s\}^{\omega} \models \mathbf{F G} p \rightarrow \mathbf{G F}(q \vee r)$ since $p$ does not occur eventually always.
- $\{q, r\} \emptyset\{p, q\} \emptyset\{r, s\}^{\omega} \models \neg(r \mathbf{U} \mathbf{X} p) \mathbf{U}(q \wedge \neg \mathbf{X X} s)$ since $q \wedge \neg \mathbf{X X} s$ already holds at the first position, i.e. $q$ occurs at the first position and $s$ does not occur at the third position.


## Solution 3.5

In the following table, $\sigma$ and $\sigma^{\prime}$ are two example sequences such that $\sigma \models \varphi$ and $\sigma^{\prime} \not \models \varphi$.

| $\varphi$ | $\sigma$ | $\sigma^{\prime}$ |
| :--- | :--- | :--- |
| (a) $\mathbf{F} g$ | $\{g\} \emptyset^{\omega}$ | $\emptyset^{\omega}$ |

(b) $\mathbf{G}(g \rightarrow \mathbf{G}(\neg s \wedge \neg r)) \quad\{g\} \emptyset^{\omega} \quad\{g, s\} \emptyset^{\omega}$ or if "after" is strict

$$
\mathbf{G}(g \rightarrow \mathbf{X G}(\neg s \wedge \neg r)) \quad\{g\} \emptyset^{\omega} \quad\{g\}\{s\} \emptyset^{\omega}
$$

(d) $\mathbf{F} g \wedge \mathbf{G}(g \rightarrow \mathbf{X G} \neg g) \quad\{g\} \emptyset^{\omega} \quad\{g\}\{g\} \emptyset^{\omega}$
(f) $\quad(\neg s \wedge \neg g) \mathbf{W} r \quad\{r\}\{g\}^{\omega} \quad\{g\}^{\omega}$

## Solution 3.6

(a) $\mathbf{G} p \rightarrow \mathbf{F} p$ is a tautology since

$$
\begin{aligned}
\mathbf{G} p \rightarrow \mathbf{F} p & \equiv \neg \mathbf{F} \neg p \rightarrow \mathbf{F} p \\
& \equiv \mathbf{F} \neg p \vee \mathbf{F} p \\
& \equiv \mathbf{F}(\neg p \vee p) \\
& \equiv \mathbf{F} \text { true } \\
& \equiv \text { true } .
\end{aligned}
$$

(b) $\mathbf{G}(p \rightarrow q) \rightarrow(\mathbf{G} p \rightarrow \mathbf{G} q)$ is a tautology. For the sake of contradiction, suppose this is not the case. There exists $\sigma$ such that

$$
\begin{align*}
& \sigma \models \mathbf{G}(p \rightarrow q), \text { and }  \tag{1}\\
& \sigma \not \models(\mathbf{G} p \rightarrow \mathbf{G} q) . \tag{2}
\end{align*}
$$

By (2), we have

$$
\begin{aligned}
& \sigma \not \models \mathbf{G} p \text {, and } \\
& \sigma \not \models \mathbf{G} q .
\end{aligned}
$$

Therefore, there exists $k \geq 0$ such that $p \in \sigma(k)$ and $q \notin \sigma(k)$ which contradicts (11).
(c) $\mathrm{FG} p \vee \mathrm{FG} \neg p$ is not a tautology since it is not satisfied by $(\{p\}\{q\})^{\omega}$.
(d) $\neg \mathbf{F} p \rightarrow \mathbf{F} \neg \mathbf{F} p$ is a tautology since $\varphi \rightarrow \mathbf{F} \varphi$ is a tautology for every formula $\varphi$.
(e) $\neg(p \mathbf{U} q) \leftrightarrow(\neg p \mathbf{U} \neg q)$ is not a tautology. Let $\sigma=\{p\}\{q\}^{\omega}$. We have $\sigma \not \vDash \neg(p \mathbf{U} q)$ and $\sigma \models \neg p \mathbf{U} \neg q$.
(f) $(\mathbf{G} p \rightarrow \mathbf{F} q) \leftrightarrow(p \mathbf{U}(p \vee q))$ is not a tautology. Let $\sigma=\emptyset\{p, q\}^{\omega}$. We have $\sigma \models \mathbf{G} p \rightarrow$ $\mathbf{F} q$ and $\sigma \nLeftarrow(p \mathbf{U}(p \vee q))$.

## Solution 3.1

Counter example: $\left\{t_{1}=r e d, t_{2}=r e d\right\}^{\omega}$
Correct formula: $\neg \mathbf{F}\left(t_{1}=\right.$ green $\wedge t_{2}=$ green $) \equiv \mathbf{G}\left(t_{1}=\right.$ red $\vee t_{2}=$ red $)$

## Solution 3.2

$$
\mathbf{G F}\left(t_{1}=\text { green }\right) \wedge \mathbf{G F}\left(t_{2}=\text { green }\right)
$$

Solution 3.3

$$
\mathbf{G}(a \wedge f \rightarrow a \wedge \mathbf{X} d) \wedge(a \mathbf{U} k)
$$

