

Model Checking – Exercise sheet 7

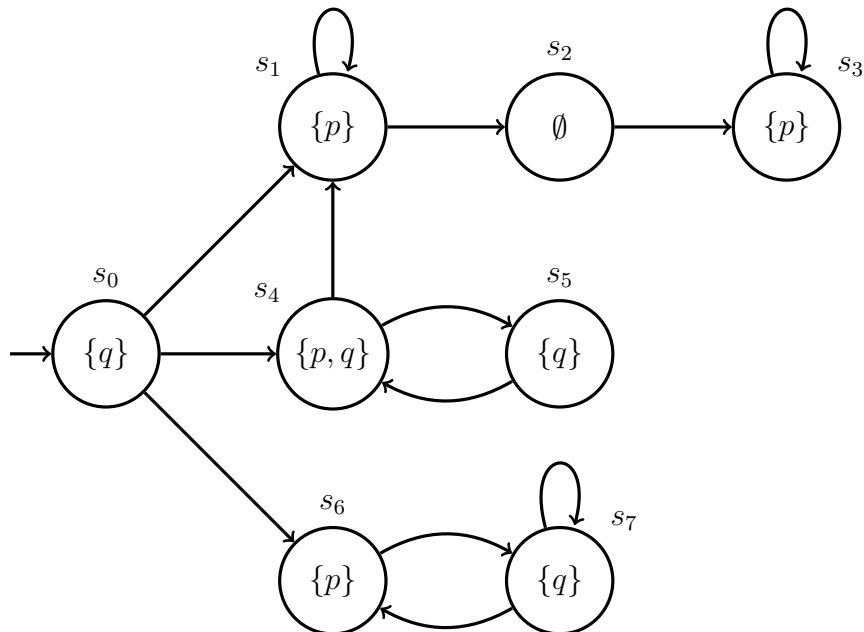
Exercise 7.1

Given two CTL formulas ϕ_1 and ϕ_2 , we write $\phi_1 \Rightarrow \phi_2$ iff for every Kripke structure \mathcal{K} we have $(\mathcal{K} \models \phi_1) \Rightarrow (\mathcal{K} \models \phi_2)$. Furthermore, we define an *implication graph* as a directed graph whose nodes are CTL formulas, and that contains an edge from ϕ_1 to ϕ_2 iff $\phi_1 \Rightarrow \phi_2$. Let $AP = \{p\}$.

- (a) Draw an implication graph with the nodes: **EFEF** p , **EGEG** p , **AFAF** p , **AGAG** p .
- (b) For each implication $\phi_1 \Rightarrow \phi_2$ obtained in (a), give a Kripke structure \mathcal{K} that satisfies ϕ_2 but not ϕ_1 , i.e. give a \mathcal{K} such that $\mathcal{K} \models \phi_2$ and $\mathcal{K} \not\models \phi_1$.
- (c) Add the following CTL formulas to the implication graph obtained in (a): **AFEF** p , **EFAF** p , **AGEG** p , **EGAG** p .
- (d) Complete the graph obtained in (c) with the nodes: **AGAF** p , **AFAG** p , **AGEF** p , **EGAF** p , **AFEG** p , **EFAG** p , **EFEG** p , **EGEF** p .

Exercise 7.2

Consider the following Kripke structure over $AP = \{p, q\}$:

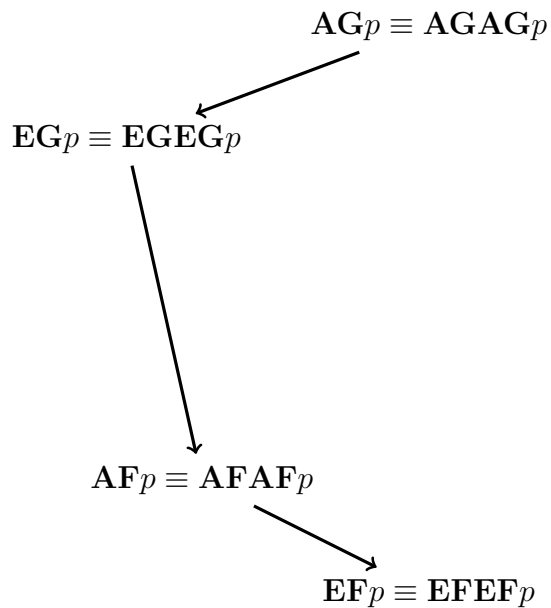


- (a) Compute $\llbracket \mathbf{EG}q \rrbracket$ and $\llbracket \mathbf{EF}q \rrbracket$.
- (b) Compute $\llbracket \mathbf{AGAF}p \rrbracket$ and $\llbracket \mathbf{EFAG}\neg q \rrbracket$.

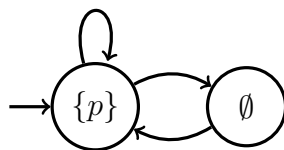
Solution 7.1

Note that the “ \Rightarrow ” relation is transitive, hence all transitive edges in (a), (b) and (d) are omitted.

(a)



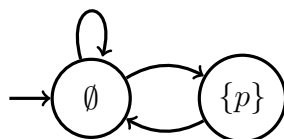
(b) The following Kripke structure satisfies $\mathbf{EG}p$, but not $\mathbf{AG}p$:



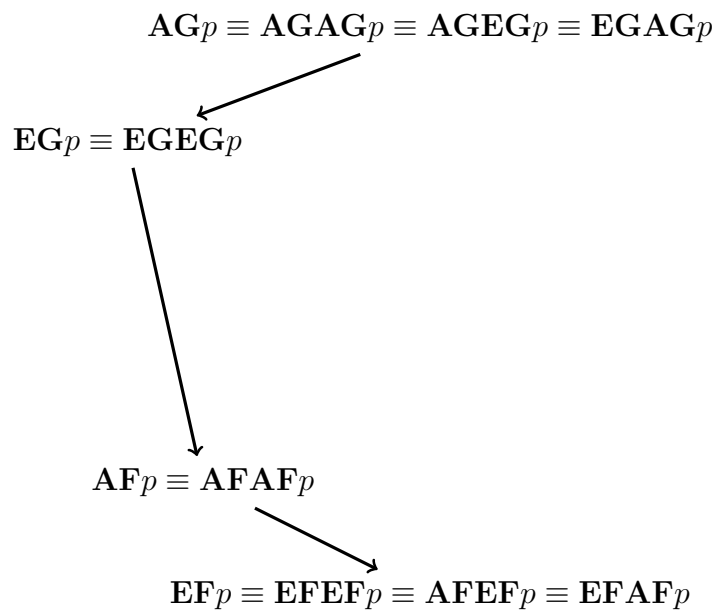
The following Kripke structure satisfies $\mathbf{AF}p$, but not $\mathbf{EG}p$:



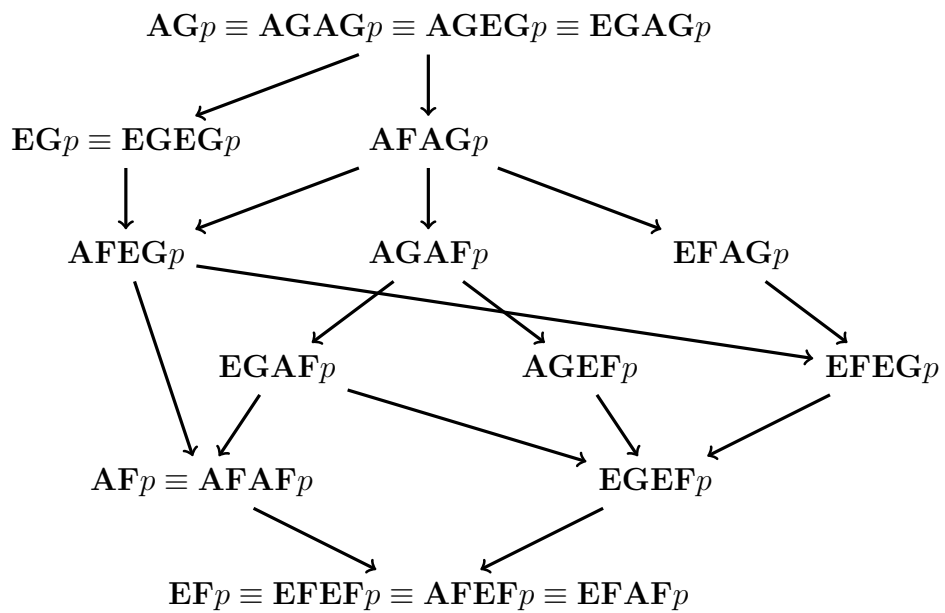
The following Kripke structure satisfies $\mathbf{EF}p$, but not $\mathbf{AF}p$:



(c)



(d)



Solution 7.2

Let $S = \{s_0, s_1, s_2, s_3, s_4, s_5, s_6, s_7\}$.

- (a) • We compute the largest fixed point from the sequence

$$\pi^0(S), \pi^1(S), \pi^2(S), \dots$$

where $\pi^0(S) = S$ and $\pi^{i+1}(S) = \llbracket q \rrbracket \cap pre(\pi^i(S))$. We obtain

$$\begin{aligned} \pi^0(S) &= S, \\ \pi^1(S) &= \{s_0, s_4, s_5, s_7\}, \\ \pi^2(S) &= \{s_0, s_4, s_5, s_7\}. \end{aligned}$$

Therefore, $\llbracket \mathbf{EG}q \rrbracket = \{s_0, s_4, s_5, s_7\}$.

- We compute the smallest fixed point from the sequence

$$\xi^0(\emptyset), \xi^1(\emptyset), \xi^2(\emptyset), \dots$$

where $\xi^0(\emptyset) = \emptyset$ and $\xi^{i+1}(\emptyset) = \llbracket q \rrbracket \cup pre(\xi^i(\emptyset))$. We obtain

$$\begin{aligned} \xi^0(\emptyset) &= \emptyset, \\ \xi^1(\emptyset) &= \{s_0, s_4, s_5, s_7\}, \\ \xi^2(\emptyset) &= \{s_0, s_4, s_5, s_6, s_7\}, \\ \xi^3(\emptyset) &= \{s_0, s_4, s_5, s_6, s_7\}. \end{aligned}$$

Therefore, $\llbracket \mathbf{EF}q \rrbracket = \{s_0, s_4, s_5, s_6, s_7\}$.

- (b) • Note that $\mathbf{AGAF}p \equiv \neg \mathbf{EFEG} \neg p$. Let us first compute $\llbracket \mathbf{EG} \neg p \rrbracket$ by computing the largest fixed point from the sequence $\pi^0(S), \pi^1(S), \pi^2(S), \dots$. We obtain

$$\begin{aligned} \pi^0(S) &= S, \\ \pi^1(S) &= \{s_0, s_2, s_5, s_7\}, \\ \pi^2(S) &= \{s_7\}, \\ \pi^3(S) &= \{s_7\}. \end{aligned}$$

Therefore, $\llbracket \mathbf{EG} \neg p \rrbracket = \{s_7\}$. In general, $\llbracket \mathbf{EF}\varphi \rrbracket$ is the set of states that can reach some state of $\llbracket \varphi \rrbracket$. By setting $\varphi = \mathbf{EG} \neg p$, we obtain

$$\begin{aligned} \llbracket \mathbf{AGAF}p \rrbracket &= \llbracket \neg \mathbf{EFEG} \neg p \rrbracket \\ &= \overline{\llbracket \mathbf{EFEG} \neg p \rrbracket} \\ &= \overline{\{s_0, s_6, s_7\}} \\ &= \{s_1, s_2, s_3, s_4, s_5\}. \end{aligned}$$

- Note that $\mathbf{EFAG}\neg q \equiv \mathbf{EF}\neg\mathbf{EF}q$. By (a), $\llbracket \mathbf{EF}q \rrbracket = \{s_0, s_4, s_5, s_6, s_7\}$, and hence $\llbracket \neg\mathbf{EF}q \rrbracket = \overline{\llbracket \mathbf{EF}q \rrbracket} = \{s_1, s_2, s_3\}$. In general, $\llbracket \mathbf{EF}\varphi \rrbracket$ is the set of states that can reach some state of $\llbracket \varphi \rrbracket$. By setting $\varphi = \neg\mathbf{EF}q$, we obtain

$$\llbracket \mathbf{EFAG}\neg q \rrbracket = \llbracket \mathbf{EF}\neg\mathbf{EF}q \rrbracket = \{s_0, s_1, s_2, s_3, s_4, s_5\}.$$