

Model Checking – Exercise sheet 2

Exercise 2.1

Let $AP = \{p, q, r, s\}$. Find out whether the following sequences satisfy the formulae $\phi = \mathbf{GF}p$ and $\psi = \mathbf{FG}(p \wedge \neg q \rightarrow (r \mathbf{U} s))$:

- \emptyset^ω
- $\{p\}^\omega$
- $\{p, q\}^\omega$
- $\{p\}\{q\}^\omega$
- $(\{p\}\{q\})^\omega$
- $(\{p\}\{p, q\})^\omega$
- $(\{p\}\{r\})^\omega$
- $(\{q\}\{s\})^\omega$
- $(\{p, s\}\{q\})^\omega$
- $(\{p\}\{q\}\{r\}\{s\})^\omega$
- $(\{p, q\}\{q\}\{r\}\{s\})^\omega$
- $(\{p, r\}\{q\}\{r\}\{s\})^\omega$
- $(\{p, r\}\{q, r\}\{r\}\{r, s\})^\omega$

Exercise 2.2

Let $AP = \{s, r, g\}$ be actions of a process: sending a message, receiving a message, and giving a result, respectively. Specify the following properties in LTL, and give example sequences that satisfy and violate the formulae.

1. The process always gives a result.
2. The process stops communicating after giving its result.
3. The process sends infinitely many messages.

4. The process only gives a result only once.
5. The process receives a message after it sends one.
6. The process does nothing until it receives a message.

Exercise 2.3

Let ϕ and ψ be LTL formulae over a set of atomic propositions AP and $\sigma \in (2^{AP})^\omega$ a valuation sequence. Prove the following equivalences:

$$\begin{aligned}
\sigma \models \mathbf{G} \phi &\Leftrightarrow \forall i : \sigma^i \models \phi \\
\sigma \models \phi \mathbf{R} \psi &\Leftrightarrow \forall i : \forall j < i : w^j \not\models \phi \rightarrow \sigma^i \models \psi \\
\sigma \models \neg(\phi \mathbf{U} \psi) &\Leftrightarrow \sigma \models \neg\phi \mathbf{R} \neg\psi \\
\sigma \models \phi \mathbf{U} \psi &\Leftrightarrow \sigma \models \psi \vee (\phi \wedge \mathbf{X}(\phi \mathbf{U} \psi)) \\
\sigma \models \phi \mathbf{R} \psi &\Leftrightarrow \sigma \models \mathbf{G} \psi \vee (\psi \mathbf{U} (\phi \wedge \psi))
\end{aligned}$$

Exercise 2.4

Let $\mathcal{K}_1 = (S, \rightarrow_1, r, AP, \nu)$ and $\mathcal{K}_2 = (S, \rightarrow_2, r, AP, \nu)$ be two almost identical Kripke structures, except for the transition relations. We write $\mathcal{K}_1 \leq \mathcal{K}_2$ if the transition relation \rightarrow_2 contains \rightarrow_1 , i.e. $\rightarrow_1 \subseteq \rightarrow_2$.

Show that if $\mathcal{K}_1 \leq \mathcal{K}_2$, then for any LTL formula ϕ :

$$\mathcal{K}_2 \models \phi \Rightarrow \mathcal{K}_1 \models \phi .$$