

## Model Checking – Solution Sheet 8

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### Exercise 8.1: What is syntactic sugar ?

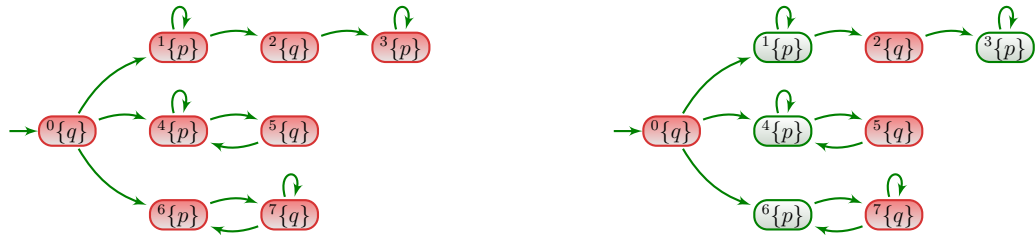
CTL operators are defined as  $QT$  where  $Q$  is  $A$  or  $E$  and  $T$  is any LTL modality ( $\mathbf{X}, \mathbf{F}, \mathbf{G}, \mathcal{U}, \mathcal{W}, \mathcal{R}$ ). So many operators means a lot of cases to handle for inductively proving results on CTL.

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|--|---|
| $\mathbf{AX}\varphi \equiv \neg\mathbf{EX}\neg\varphi$ $\mathbf{EF}\varphi \equiv \text{true } \mathbf{EU}\varphi$ <ol style="list-style-type: none"> <li>1. <math>\mathbf{AG}\varphi \equiv \neg\mathbf{EF}\neg\varphi</math></li> <li style="padding-left: 20px;"><math>\mathbf{AF}\varphi \equiv \neg\mathbf{EG}\neg\varphi</math></li> </ol> | $\varphi_1 \mathbf{E}\mathcal{W}\varphi_2 \equiv \mathbf{EG}\varphi_1 \vee (\varphi_1 \mathbf{EU}\varphi_2)$ $\varphi_1 \mathbf{A}\mathcal{W}\varphi_2 \equiv \neg(\neg\varphi_2 \mathbf{EU}\neg(\varphi_1 \vee \varphi_2))$ $\varphi_1 \mathbf{A}\mathcal{U}\varphi_2 \equiv \mathbf{AF}\varphi_2 \wedge (\varphi_1 \mathbf{A}\mathcal{W}\varphi_2)$ $\varphi_1 \mathbf{E}\mathcal{R}\varphi_2 \equiv \neg(\neg\varphi_1 \mathbf{A}\mathcal{U}\neg\varphi_2)$ $\varphi_1 \mathbf{A}\mathcal{R}\varphi_2 \equiv \neg(\neg\varphi_1 \mathbf{EU}\neg\varphi_2)$ |
|--|---|

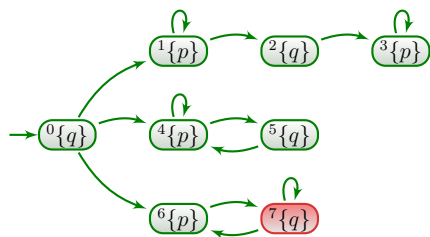
2.  $\mathbf{EX}$  only provides a local information on the direct successors.  $\mathbf{EU}$  only provides the existence of a branch which admits a treshold: the converse is the absence of a treshold on all branch.  $\mathbf{EG}$  is thus necessary as it allows to state the absence of a treshold on a branch (or the presence of one on every branches).

### Exercise 8.3: Fixpoint computations

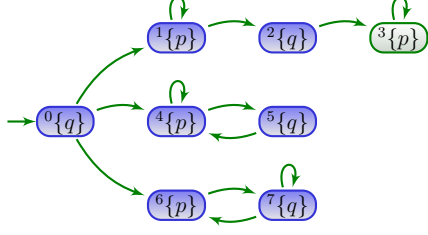
1. We recall that  $\mathbf{EG}q$  is computed as the greatest fixpoint of the equation  $X = \mu(q) \cap \text{pre}(X)$ . Thus we start by considering  $X = [0, 7]$ .



After the first iteration, since  $\text{pre}([0, 7]) = [0, 6]$ , we have all states at which  $q$  holds. Only 7 then has a predecessor for which  $q$  holds. The next iteration does not modify the set  $X$  computed, therefore follows states that validate  $\mathbf{EG}q$ .



$EFq$  is computed as least fix point of the equation  $X = \mu(q) \cup (pre(X))$ , ie. initially  $X = \emptyset$ . After 3 iteration, we obtain the following fixpoint:



2. Remark that  $AGAFp = \neg EFEG\neg p$ . We start by computing  $\llbracket EG\neg p \rrbracket$ . Remark that  $\llbracket \neg p \rrbracket = \llbracket q \rrbracket$ . Thus  $\llbracket EG\neg p \rrbracket = \llbracket EGq \rrbracket = \{7\}$ .

Next we compute  $\llbracket EF\{7\} \rrbracket$ , as the least fixpoint of  $X = \mu(q) \cup (pre(X))$ .

$$X_0 = \emptyset, pre(X_0) = \emptyset, \mu(\{7\}) = \{7\}$$

$$X_1 = \{7\}, pre(\{7\}) = \{7, 6\}, \mu(\{7\}) = \{7\}$$

$$X_2 = \{7, 6\}, pre(\{7, 6\}) = \{7, 6, 0\}, \mu(\{7\}) = \{7\}$$

$$X_3 = \{7, 6, 0\}, pre(X_3) = X_3$$

$$\llbracket EF\{7\} \rrbracket = X_3 = \{0, 6, 7\}$$

$$\llbracket \neg EF\{7\} \rrbracket = \mathbb{C}\llbracket \neg EF\{7\} \rrbracket = \{1, 2, 3, 4, 5\}$$

$$\text{Therefore } \llbracket AGAFp \rrbracket = \{1, 2, 3, 4, 5\}$$

Remark that  $AFAGp = \neg EGEF\neg p$ . We start by computing  $\llbracket EF\neg p \rrbracket$ . Remark that  $\llbracket \neg p \rrbracket = \llbracket q \rrbracket$ . Thus  $\llbracket EF\neg p \rrbracket = \llbracket EFq \rrbracket = \{0, 1, 2, 4, 5, 6, 7\}$ .

Next we compute  $\llbracket EG\{0, 1, 2, 4, 5, 6, 7\} \rrbracket$ , as the greatest fixpoint of the equation  $X = \mu(\{0, 1, 2, 4, 5, 6, 7\}) \cap pre(X)$

$$X_0 = [0, 7], pre(X_0) = X_0, \mu(\{0, 1, 2, 4, 5, 6, 7\}) = \{0, 1, 2, 4, 5, 6, 7\}, \text{ thus}$$

$$X_1 = \{0, 1, 2, 4, 5, 6, 7\}, pre(X_1) = \{0, 1, 4, 5, 6, 7\}, \text{ therefore}$$

$$X_2 = \{0, 1, 4, 5, 6, 7\}, pre(X_2) = \{0, 1, 4, 5, 6, 7\},$$

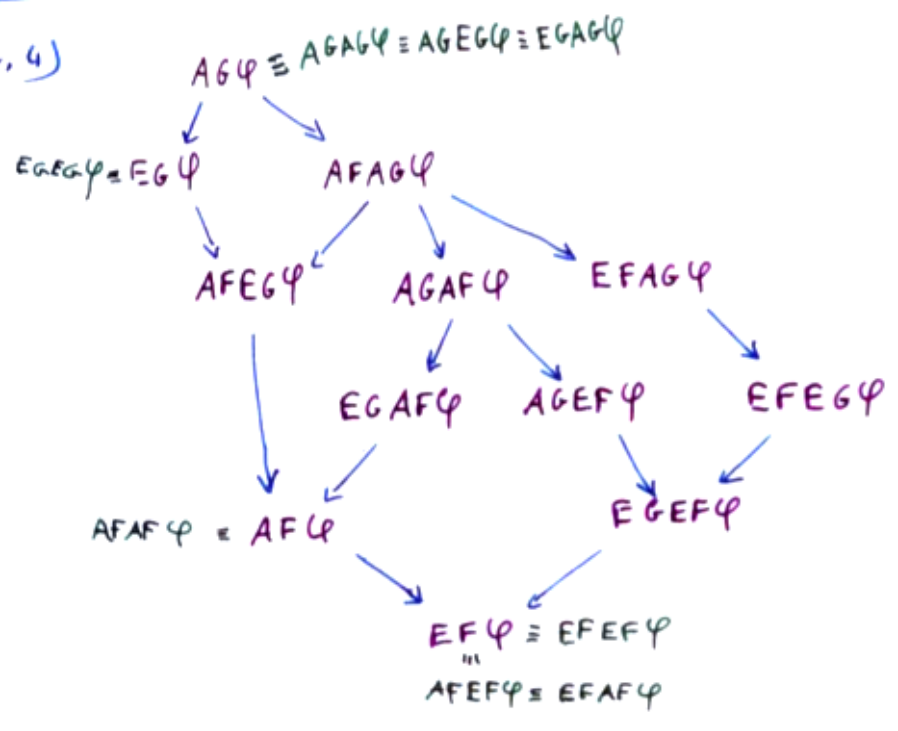
$$\text{so } \llbracket EG\{0, 1, 2, 4, 5, 6, 7\} \rrbracket = X_2 = \{0, 1, 4, 5, 6, 7\}. \text{ Therefore } \llbracket \neg EG\{0, 1, 2, 4, 5, 6, 7\} \rrbracket = \mathbb{C}\{0, 1, 4, 5, 6, 7\} = \{2, 3\}.$$

$$\llbracket AFAGp \rrbracket = \{2, 3\}$$

3. Remark the run  $0, 6, 7, 7, \dots$ . Its trace is  $\{q\}\{p\}\{q\}^\omega$  which validates neither  $\mathbf{FG}p$ , nor  $\mathbf{GF}p$ . Thus  $K \not\models \mathbf{FG}p$  and  $K \not\models \mathbf{GF}p$ .

Ex 8.2

1, 2, 4)



3)

