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# Model Checking – Exercise sheet 4

## Exercise 4.1

We give the following regular expressions over defining regular languages over the alphabet  $\Sigma = \{a, b\}$  (*C* denotes the complement w.r.t.  $\Sigma^*$ ).

 $L_1 = a^*(baa^*)ba^*$   $L_2 = (a^*ba)^*a^*ba^*$   $L_3 = C(\Sigma^*bb\Sigma^*)$ 

- 1. Which of those regular expressions denote the same langage ?
- 2. Which is (are) the regular expression(s) accepted by the smallest deterministic automaton ?

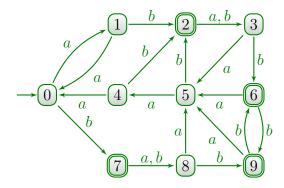
#### Exercise 4.2

Give a regular expression for each of the following languages over alphabet  $\Sigma = \{a, b\}$ :

- 1. words that start with b and don't contain pattern ab
- 2. words in which any two b are always separated by at least one a
- 3. words that do not contain the pattern *aab*
- 4. words containing either two a or three b
- 5. words with as many a than b
- 6. words in which at each position the number of preceeding a is the same as the number of preceeding b up to a difference of 3.
- 7. words in which the pattern ab occurs as often as the pattern ba

## **Exercise 4.3:** Minimization

Minimize the following automaton:



#### Exercise 4.4: True or False

Prove or disprove each of the following assertions:

- 1. Let  $A_1 = (q_0, Q_1, \delta_1, F_1)$  and  $A_2 = (q'_0, Q_2, \delta_2, F_2)$  two deterministic finite automata. The minimal automaton accepting  $\mathcal{L}(A_1) \cup \mathcal{L}(A_2)$  may have  $|Q_1| \cdot |Q_2|$  states.
- 2. For any non-deterministic automaton, one can build a deterministic automaton accepting the same language.
- 3. The minimal deterministic automaton of a non deterministic automaton  $A = (Q_0, Q, \Delta, F)$  may have a number of states exponential w.r.t. Q.
- 4. The reverse language of a regular language is regular.
- 5. It suffices to reverse the transition and change accepting for initial and initial to accepting on an automaton to accept the reverse of the language accepted by an automaton.
- 6. To test whether two automata accept the same language we can test whether the two reverse automata accept the same language.
- 7. To test whether two deterministic automata accept the same language, we can test whether determizing the two reverse automata lead to the same automaton.
- 8. The complement of a regular language is a regular language. It suffices make every accepting state non-accepting and vice-versa to build an automaton accepting the complement of that language.
- 9. Any regular language is accepted by a minimal deterministic automaton.
- 10. Any regular language is accepted by a minimal non-deterministic automaton.
- 11. Any regular language admits a minimal regular expression.

## Exercise 4.5: And for infinite words

Prove or disprove propositions 1,2,3,8,9,10 for infinite words.