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## Logic

## Exercise Sheet 9

Discussion: July 10, 2014

- 1. Give a family of formulas  $(F_k)_{k\geq 0}$  of predicate logic over the signature  $S = \{E\}$ , where E is a binary symbol such that the following holds for each  $k \geq 0$ :
  - $\operatorname{qr}(F_k) = k$ ,
  - $F_k = F_k(x, y)$ , i.e.  $F_k$  has two free variables, and
  - for each suitable structure  $\mathcal{A}$  and each  $a, b \in U_{\mathcal{A}}$  we have  $\mathcal{A}_{[x/a,y/b]} \models F_k$  if and only if there is *E*-path from *a* to *b* of length *k*.

Construct a similar family  $(F_k)_{k\geq 0}$  but where

- for each suitable structure  $\mathcal{A}$  and each  $a, b \in U_{\mathcal{A}}$  we have  $\mathcal{A}_{[x/a,y/b]} \models F_k$  if and only if there is *E*-path from *a* to *b* of length  $2^k$ .
- 2. Prove the second part of the induction proof for showing connectivity is not expressible in predicate logic. More precisely, consider the case  $a \notin N_{2^{k-(i+1)}}(a_h)$  for each  $h \in \{1, \ldots, i\}$ .
- **3.** Show that there is no formula in predicate logic expressing the transitive closure of a binary relation R.
- **4.** Show the following equivalence for each subset  $X \subseteq \mathbb{N}$  and each signature S:
  - (1) Property  $\{\mathcal{A} : |U_{\mathcal{A}}| \in X\}$  is expressible in predicate logic with equality over signature S.
  - (2) X or  $\mathbb{N} \setminus X$  is finite.