

4 Procedures as triples

(fun f M = e, t, V)

5 Evaluation

$V \models e : v$ denotes that in the value environment V the expression e evaluates to the value v .

$V \models p : V1$ denotes that in the value environment V the evaluation of the program p results in the value environment $V1$.

5.1 Expressions

$V(b) = v$

 $V \models b : v \quad V \models k : k$

$V \models e : v \quad V \models e1 : v1 \quad \dots \quad V \models eN : vN \quad V \models e : (v1, \dots, vN)$

 $V \models (e) : v \quad V \models (e1, \dots, eN) : (v1, \dots, vN) \quad V \models \#Ie : vI$

$V \models e : v \quad V \models e1 : v1 \quad V \models e2 : v2$

 $V \models o e : o v \quad V \models e1 o e2 : v1 o v2$

$V \models e1 : \text{true} \quad V \models e2 : v \quad V \models e1 : \text{false} \quad V \models e3 : v$

 $V \models \text{if } e1 \text{ then } e2 \text{ else } e3 : v \quad V \models \text{if } e1 \text{ then } e2 \text{ else } e3 : v$

$V \models e1 : (\text{fun } f \text{ } b = e, t, V1) \quad V \models e2 : v2$
 $V1 + [f := (\text{fun } f \text{ } b = e, t, V1)] + [b := v2] \models e : v$

 $V \models e1 e2 : v$

$V \models p : V1 \quad V1 \models e : v$

 $V \models \text{let } p \text{ in } e \text{ end} : v$

5.2 Declarations

$V \models e : v$

 $V \models \text{val } b = e : V + [b := v]$

$V1 = (V \text{ restricted to } \text{FreeIds}(\text{fun } f (b : t1) : t2 = e))$

 $V \models \text{fun } f (b : t1) : t2 = e : V + [f := (\text{fun } f \text{ } b = e, t1 \rightarrow t2, V1)]$

5.3 Programs

$V0 \models d1 : V1 \quad \dots \quad V(N-1) \models dN : VN$

 $V0 \models d1 \dots dN : VN$

5.4 Examples

$[x:=1] \models x : 1 \quad [x:=1] \models 3 : 3$

 $[x:=1] \models x+3 : 4$

$[x:=1, a:=2, f:= (\text{fun } f \text{ } x = x+a, \text{int} \rightarrow \text{int}, [a:=2])] \models f : (\text{fun } f \text{ } x = x+a, \text{int} \rightarrow \text{int}, [a:=2])$

$[x:=1, a:=2, f:= (\text{fun } f \text{ } x = x+a, \text{int} \rightarrow \text{int}, [a:=2])] \models x+3 : 4$

$[a:=2] + [f : (\text{fun } f \text{ } x = x+1, \text{int} \rightarrow \text{int}, [a:=2])] + [x:=4] \models x+a : 5$

 $[x:=1, a:=2, f:= (\text{fun } f \text{ } x = x+a, \text{int} \rightarrow \text{int}, [a:=2])] \models f (x+3) : 6$

```
[] |>> val x = 1 : [x:=1]
[x:=1] |>> val a = 2 : [x:=1, a:=2]
[x:=1, a:=2] |>> fun f (x:int):int = x+a : [x:=1, a:=2, f:= (fun f x = x+a, int -> int, [a:=2])]
[x:=1, a:=2, f:= (fun f x = x+a, int -> int, [a:=2])] |>> val y = f (x+3) : ... + [y:=6]
-----
[] |>> val x = 1 val a = 2 fun f (x:int):int = x+a val y = f (x+3)
      : [x:=1, a:=2, f:= (fun f x = x+a, int -> int, [a:=2]), y:=6]
```

6 Procedures w/o names

```
2;  
  
it + it;  
  
let fun f x = x+1 in f end;  
  
it 1;  
  
val f = let fun f (x:int) = x+1 in f end;  
  
f 1;  
  
fn x => x+1;  
  
it 1;  
  
val f = (fn x => x+1);  
  
f 1;  
  
val f = fn x => x+1;  
  
f 1;
```

7 Curried (kaskadierte)procedures

```
val add =  
  (fn x =>  
    (fn y =>  
      x+y  
    )  
  );  
  
add 1 2;  
  
fun add x y = x+y  
  
add 1;  
  
it 2;  
  
val inc = add 1;  
  
inc 1;  
  
(add 1) 2  
  
int -> (int -> int)
```

8 Iterators

```
sum_seq n = 0 + 1 + ... + n

fun sum_seq (n:int) = if n < 1 then 0 else n + sum_seq (n-1)

sum_of_sqr n = 0 + 1*1 + ... + n*n

fun sum_sqr (n:int) = if n < 1 then 0 else n*n + sum_sqr (n-1)

sum f n = 0 + f 1 + ... + f n

sum_seq n = sum (fn x => x) n
sum_sqr n = sum (fn x => x*x) n

sum: (int->int) -> int -> int

fun sum (f : int->int) (n:int) = if n < 1 then 0 else sum f (n-1) + f n

sum_seq 3 = sum (fn x => x) 3

sum_of_squares 3 = sum (fn x => x*x) 3

% bounded iteration

iter n s f = f(...(f s)...)

iter 3 s f = f(f(f(s)))

fun iter (n:int) (s:int) (f:int->int) = if n<1 then s else iter (n-1) (f s) f

fun two_power_n n = iter n 1 (fn x => 2*x)

% unbounded iteration

first s p = minimal t >= s such that p s is true

first 0 (fn x => x >= 3) = 3

fun first (s:int) (p : int -> bool) = if p s then s else first (s+1) p
```