

# 1 Syntax

$\langle \text{Aus} \rangle ::= \langle \text{Konst} \rangle \mid \langle \text{Bez} \rangle \mid (\langle \text{Aus} \rangle) \mid (\langle \text{Aus} \rangle, \dots, \langle \text{Aus} \rangle) \mid \# \langle \text{ganze pos. Zahl} \rangle \langle \text{Aus} \rangle$   
 $\mid \langle \text{einst. Op} \rangle \langle \text{Aus} \rangle \mid \langle \text{Aus} \rangle \langle \text{zweist. Op} \rangle \langle \text{Aus} \rangle$   
 $\mid \text{if } \langle \text{Aus} \rangle \text{ then } \langle \text{Aus} \rangle \text{ else } \langle \text{Aus} \rangle \mid \langle \text{Aus} \rangle \langle \text{Aus} \rangle \mid \text{let } \langle \text{Prog} \rangle \text{ in } \langle \text{Aus} \rangle \text{ end}$

$\langle \text{Typ} \rangle ::= \text{int} \mid \text{real} \mid \text{bool} \mid \text{unit} \mid (\langle \text{Typ} \rangle) \mid (\langle \text{Typ} \rangle * \dots * \langle \text{Typ} \rangle) \mid \langle \text{Typ} \rangle \rightarrow \langle \text{Typ} \rangle$

$\langle \text{Dek} \rangle ::= \text{val } \langle \text{Val-Muster} \rangle = \langle \text{Aus} \rangle$   
 $\mid \text{fun } \langle \text{Bez} \rangle \langle \text{Arg-Muster} \rangle = \langle \text{Aus} \rangle \mid \text{fun } \langle \text{Bez} \rangle \langle \text{Arg-Muster} \rangle : \langle \text{Typ} \rangle = \langle \text{Aus} \rangle$

$\langle \text{Val-Muster} \rangle ::= \langle \text{Bez} \rangle \mid (\langle \text{Bez} \rangle, \dots, \langle \text{Bez} \rangle)$

$\langle \text{Arg-Muster} \rangle ::= \langle \text{Bez} \rangle : \langle \text{Typ} \rangle \mid (\langle \text{Bez} \rangle : \langle \text{Typ} \rangle, \dots, \langle \text{Bez} \rangle : \langle \text{Typ} \rangle)$

$\langle \text{Prog} \rangle ::= \langle \text{Dek} \rangle \dots \langle \text{Dek} \rangle$

## 2 Environment

Function from identifiers to types or values. Composition of environments  $f + g$  overwrites bindings in  $f$ .

$$(f + g)(x) = \begin{cases} g(x), & \text{if } x \in \text{dom}(g) \\ f(x), & \text{otherwise} \end{cases}$$

## 3 Typing

$T \mid - e : t$  denotes that in the type environment  $T$  the expression  $e$  has the type  $t$ .

$T \mid > p : T1$  denotes that in the type environment  $T$  the typing of the program  $p$  results in the type environment  $T1$ .

### 3.1 Expressions

$T(b) = t$	$k \in \{\text{true}, \text{false}\}$	$k \in \mathbb{Z}$	$k \in \mathbb{R}$
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$T \mid - b : t$	$T \mid - k : \text{bool}$	$T \mid - k : \text{int}$	$T \mid - k : \text{real} \quad T \mid - () : \text{unit}$
$T \mid - e : t$	$T \mid - e1 : t1 \dots T \mid - eN : tN$		$T \mid - e : t1 * \dots * tN \quad I \in 1..N$
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$T \mid - (e) : t$	$T \mid - (e1, \dots, eN) : t1 * \dots * tN$		$T \mid - \#Ie : tI$
$T \mid - o : t1 \rightarrow t2$	$T \mid - e : t1$	$T \mid - e1 : t1$	$T \mid - o : t1 * t2 \rightarrow t \quad T \mid - e2 : t2$
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$T \mid - oe : t2$		$T \mid - e1 o e2 : t$	
$T \mid - e1 : \text{bool}$	$T \mid - e2 : t$	$T \mid - e3 : t$	$T \mid - e1 : t1 \rightarrow t2 \quad T \mid - e2 : t1$
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$T \mid - \text{if } e1 \text{ then } e2 \text{ else } e3 : t$			$T \mid - e1 e2 : t2$
$T \mid > p : T1$	$T1 \mid - e : t$		
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$T \mid - \text{let } p \text{ in } e \text{ end} : t$			

### 3.2 Declarations

$T \mid - e : t$	$T + [f := t1 \rightarrow t2] + [b := t1] \mid - e : t2$
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$T \mid > \text{val } b = e : T + [b := t]$	$T \mid > \text{fun } f (b : t1) : t2 = e : T + [f := t1 \rightarrow t2]$

### 3.3 Programs

$T0 \mid > d1 : T1 \dots TN \mid > dN : T(N+1)$   
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 $T0 \mid > d1 \dots dN : T(N+1)$