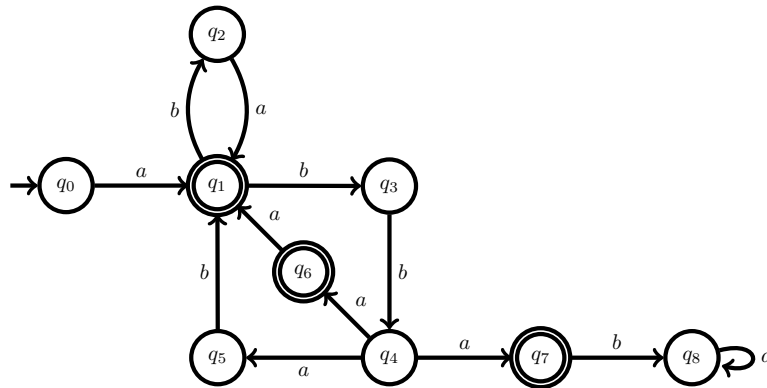


## Automata and Formal Languages — Homework 13

Due 29.01.2019

### Exercise 13.1

Let  $B$  be the following Büchi automaton:



- (a) Execute the emptiness algorithm *NestedDFS* on  $B$ .
- (b) Recall that *NestedDFS* is a non-deterministic algorithm and different choices of runs may return different lassos. Which lassos of  $B$  can be found by *NestedDFS*?
- (c) Show that *NestedDFS* is non optimal by exhibiting some search sequence on  $B$ .
- (d) Execute the emptiness algorithm *TwoStack* on  $B$ .
- (e) Which lassos of  $B$  can be found by *TwoStack*?

### Exercise 13.2

Compare the following two formulas using Spot (<https://spot.lrde.epita.fr/app/>). Determine whether the languages of the formulas are equal or not. In the case they are not equal give a distinguishing word.

1. After  $q$  until  $r$ :

$$\mathbf{G}(q \rightarrow ((\neg p \wedge \neg r) \mathbf{U} (r \vee ((p \wedge \neg r) \mathbf{U} (r \vee ((\neg p \wedge \neg r) \mathbf{U} (r \vee ((p \wedge \neg r) \mathbf{U} (r \vee (\neg p \mathbf{W} r) \vee \mathbf{G}p))))))))))$$

2. Between  $q$  and  $r$ :

$$\mathbf{G}((q \wedge \mathbf{F}r) \rightarrow ((\neg p \wedge \neg r) \mathbf{U} (r \vee ((p \wedge \neg r) \mathbf{U} (r \vee ((\neg p \wedge \neg r) \mathbf{U} (r \vee ((p \wedge \neg r) \mathbf{U} (r \vee (\neg p \mathbf{U} r))))))))))$$

**Exercise 13.3**

Prove or disprove:

(a)  $\mathbf{GF}(\varphi \vee \psi) \equiv \mathbf{GF}\varphi \vee \mathbf{GF}\psi$

(b)  $\mathbf{GF}(\varphi \wedge \psi) \equiv \mathbf{GF}\varphi \wedge \mathbf{GF}\psi$

(c)  $(\varphi \vee \psi) \mathbf{U} \rho \equiv (\varphi \mathbf{U} \rho) \vee (\psi \mathbf{U} \rho)$

(d)  $\rho \mathbf{U} (\varphi \vee \psi) \equiv (\rho \mathbf{U} \varphi) \vee (\rho \mathbf{U} \psi)$

(e)  $\mathbf{FGX}\varphi \equiv \mathbf{FG}\varphi$

(f)  $(\mathbf{GF}\varphi) \wedge (\mathbf{GF}\psi) \equiv \mathbf{GF}(\varphi \wedge \mathbf{F}\psi)$

**Exercise 13.4**

Let  $\text{AP} = \{p, q, r\}$ . Give formulas for the computations satisfying the following properties:

(a) if  $q$  eventually holds, then  $p$  may not hold before  $q$  first do.

(b) if  $q$  eventually holds, then  $p$  holds before  $q$  first holds.

(c)  $p$  always holds between  $q$  and  $r$ .

(d)  $p$ , and only  $p$ , holds at even positions and  $q$ , and only  $q$ , holds at odd positions.

### Solution 13.1

- (a) Let us assume that the algorithms always pick states in ascending order with respect to their indices. *dfs1* visits  $q_0, q_1, q_2, q_3, q_4, q_5, q_6$ , then calls *dfs2* which visits  $q_6, q_1, q_2, q_3, q_4, q_5, q_6$  and reports “non empty”.
- (b) Since  $q_7$  does not belong to any lasso, only lassos containing  $q_1$  or  $q_6$  can be found. In every run of the algorithm, *dfs1* blackens  $q_6$  before  $q_1$ . The only lasso containing  $q_6$  is:  $q_0, q_1, q_3, q_4, q_6, q_1$ . Therefore, this is the only lasso that can be found by the algorithm.
- (c) The execution given in (a) shows that *NestedDFS* is non optimal since it returns the lasso  $q_0, q_1, q_3, q_4, q_6, q_1$  even though the lasso  $q_0, q_1, q_2, q_1$  was already appearing in the explored subgraph.
- (d) Let us assume that the algorithms always pick states in ascending order with respect to their indices. The algorithm reports “non empty” after the following execution:

$\xrightarrow{\begin{array}{l} C.\text{push}(q_0) \\ V.\text{push}(q_0) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_1) \\ V.\text{push}(q_1) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_2) \\ V.\text{push}(q_2) \end{array}}$	$C$	$V$	$\xrightarrow{C.\text{pop}()}$	$C$	$V$	$\xrightarrow{C.\text{pop}()}$	$C$	$V$
	$q_0$	$q_0$		$q_1$	$q_1$		$q_2$	$q_2$		$q_1$	$q_1$		$q_0$	$q_0$
				$q_0$	$q_0$		$q_1$	$q_1$		$q_1$	$q_1$		$q_0$	$q_0$

- (e) All of them. The lasso  $q_0, q_1, q_2, q_1$  is found by the above execution. The lasso  $q_0, q_1, q_3, q_4, q_6, q_1$  is found by the following execution:

$\xrightarrow{\begin{array}{l} C.\text{push}(q_0) \\ V.\text{push}(q_0) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_1) \\ V.\text{push}(q_1) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_3) \\ V.\text{push}(q_3) \end{array}}$	$C$	$V$
	$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$
				$q_0$	$q_0$		$q_1$	$q_1$
							$q_0$	$q_0$
							$q_3$	$q_3$
							$q_1$	$q_1$
							$q_0$	$q_0$
							$q_4$	$q_4$
							$q_3$	$q_3$
							$q_1$	$q_1$
							$q_0$	$q_0$
							$q_6$	$q_6$
							$q_4$	$q_4$
							$q_3$	$q_3$
							$q_1$	$q_1$
							$q_0$	$q_0$
							$q_6$	$q_6$
							$q_4$	$q_4$
							$q_3$	$q_3$
							$q_1$	$q_1$
							$q_0$	$q_0$

The lasso  $q_0, q_1, q_3, q_4, q_5, q_1$  is found by the following execution:

$\xrightarrow{\begin{array}{l} C.\text{push}(q_0) \\ V.\text{push}(q_0) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_1) \\ V.\text{push}(q_1) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_3) \\ V.\text{push}(q_3) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_4) \\ V.\text{push}(q_4) \end{array}}$	$C$	$V$	$\xrightarrow{\begin{array}{l} C.\text{push}(q_5) \\ V.\text{push}(q_5) \end{array}}$	$C$	$V$
	$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$		$q_4$	$q_4$		$q_5$	$q_5$
				$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$		$q_4$	$q_4$
							$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$
							$q_0$	$q_0$		$q_3$	$q_3$		$q_4$	$q_4$
							$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$
							$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$
							$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$
							$q_0$	$q_0$		$q_1$	$q_1$		$q_3$	$q_3$

### Solution 13.3

- (a) True. If  $\sigma \models \mathbf{GF}\varphi \vee \mathbf{GF}\psi$ , then  $\sigma \models \mathbf{GF}(\varphi \vee \psi)$ . If  $\sigma \models \mathbf{GF}(\varphi \vee \psi)$ , then there exist  $i_0 < i_1 < \dots$  such that

$$\sigma^{i_j} \models \varphi \vee \psi \text{ for every } j \in \mathbb{N}. \quad (1)$$

Let  $I = \{j \in \mathbb{N} : \sigma^{i_j} \models \varphi\}$  and  $J = \{j \in \mathbb{N} : \sigma^{i_j} \models \psi\}$ . If  $I$  and  $J$  are both finite, then (1) does not hold, which is a contradiction. Therefore, at least one of  $I$  and  $J$  is infinite. This implies that  $\sigma \models \mathbf{GF}\varphi \vee \mathbf{GF}\psi$ .  $\square$

- (b) False. Let  $\sigma = (\{p\}\{q\})^\omega$ . We have  $\sigma \not\models \mathbf{GF}(p \wedge q)$  and  $\sigma \models \mathbf{GF}p \wedge \mathbf{GF}q$ .
- (c) False. Let  $\sigma = \{p\}\{q\}\{r\}\emptyset^\omega$ . We have  $\sigma \models (p \vee q) \mathbf{U} r$  and  $\sigma \not\models (p \mathbf{U} r) \vee (q \mathbf{U} r)$ .

(d) True, since:

$$\begin{aligned}\sigma \models \rho \mathbf{U} (\varphi \vee \psi) &\iff \exists k \geq 0 \text{ s.t. } \sigma^k \models (\varphi \vee \psi) \wedge \forall 0 \leq i < k \sigma^i \models \rho \\ &\iff \exists k \geq 0 \text{ s.t. } ((\sigma^k \models \varphi) \vee (\sigma^k \models \psi)) \wedge \forall 0 \leq i < k \sigma^i \models \rho \\ &\iff \exists k \geq 0 \text{ s.t. } (\sigma^k \models \varphi \wedge \forall 0 \leq i < k \sigma^i \models \rho) \vee (\sigma^k \models \psi \wedge \forall 0 \leq i < k \sigma^i \models \rho) \\ &\iff (\exists k \geq 0 \text{ s.t. } \sigma^k \models \varphi \wedge \forall 0 \leq i < k \sigma^i \models \rho) \vee (\exists k \geq 0 \text{ s.t. } \sigma^k \models \psi \text{ and } \forall 0 \leq i < k \sigma^i \models \rho) \\ &\iff \sigma \models (\rho \mathbf{U} \varphi) \vee (\rho \mathbf{U} \psi). \quad \square\end{aligned}$$

#### Solution 13.4

- (a)  $\mathbf{F}q \rightarrow (\neg p \mathbf{U} q)$
- (b)  $\mathbf{F}q \rightarrow (\neg q \mathbf{U} (\neg q \wedge p))$
- (c)  $\mathbf{G}((q \wedge \mathbf{X}\mathbf{F}r) \rightarrow \mathbf{X}(p \mathbf{U} r))$
- (d)  $\mathbf{G}(\neg r) \wedge \mathbf{G}(p \leftrightarrow \neg q) \wedge p \wedge \mathbf{G}(p \rightarrow \mathbf{X}q) \wedge \mathbf{G}(q \rightarrow \mathbf{X}p)$