

Automata and Formal Languages — Homework 11

Due 16.01.2018

Exercise 11.1

Let $\text{inf}(w)$ denote the set of letters occurring infinitely often in the infinite word w . Give Büchi automata and ω -regular expressions for the following ω -languages over $\Sigma = \{a, b, c\}$:

- (a) $L_1 = \{w \in \Sigma^\omega : \text{inf}(w) \subseteq \{a, b\}\}$,
- (b) $L_2 = \{w \in \Sigma^\omega : \text{inf}(w) = \{a, b\}\}$,
- (c) $L_3 = \{w \in \Sigma^\omega : \{a, b\} \subseteq \text{inf}(w)\}$,
- (d) $L_4 = \{w \in \Sigma^\omega : \text{inf}(w) = \{a, b, c\}\}$.

Exercise 11.2

Give *deterministic* Büchi automata recognizing the following ω -languages over $\Sigma = \{a, b, c\}$:

- (a) $L_1 = \{w \in \Sigma^\omega : w \text{ contains at least one } c\}$,
- (b) $L_2 = \{w \in \Sigma^\omega : \text{in } w, \text{ every } a \text{ is immediately followed by a } b\}$,
- (c) $L_3 = \{w \in \Sigma^\omega : \text{in } w, \text{ between two successive } a\text{'s there are at least two } b\text{'s}\}$.

Exercise 11.3

Give *deterministic* Rabin automata, Muller automata and parity automata for the following language:

$$L = \{w \in \{a, b\}^\omega : w \text{ contains finitely many } a\text{'s}\}.$$

Exercise 11.4

Prove or disprove:

- (a) For every Büchi automaton A , there exists a Büchi automaton B with a single initial state and such that $L_\omega(A) = L_\omega(B)$;
- (b) For every Büchi automaton A , there exists a Büchi automaton B with a single accepting state and such that $L_\omega(A) = L_\omega(B)$;
- (c) There exists a Büchi automaton recognizing the finite ω -language $\{w\}$ such that $w \in \{0, 1, \dots, 9\}^\omega$ and w_i is the i^{th} decimal of π .