Automata and Formal Languages — Homework 8

Due 12.12.2017

Exercise 8.1
Let $L_1 = \{abb, bba, bbb\}$ and $L_2 = \{aba, bbb\}$.

(a) Give an algorithm for the following operation:

INPUT: A fixed-length language $L \subseteq \Sigma^k$ described explicitely by a set of words.
OUTPUT: State q of the master automaton over Σ such that $L(q) = L$.

(b) Use the previous algorithm to build the states of the master automaton for L_1 and L_2.

(c) Compute the state of the master automaton representing $L_1 \cup L_2$.

(d) Identify the kernels $\langle L_1 \rangle$, $\langle L_2 \rangle$, and $\langle L_1 \cup L_2 \rangle$.

Exercise 8.2

(a) Give an algorithm for the following operation:

INPUT: States p and q of the master automaton.
OUTPUT: State r of the master automaton such that $L(r) = L(p) \cdot L(q)$.

(b) A *coding* over an alphabet Σ is a function $h \colon \Sigma \mapsto \Sigma$. A coding h can naturally be extended to a morphism over words, i.e. $h(\varepsilon) = \varepsilon$ and $h(w) = h(w_1)h(w_2)\cdots h(w_n)$ for every $w \in \Sigma^n$. Give an algorithm for the following operation:

INPUT: A state q of the master automaton and a coding h.
OUTPUT: State r of the master automaton such that $L(r) = \{h(w) : w \in L(q)\}$.

Can you make your algorithm more efficient when h is a permutation?

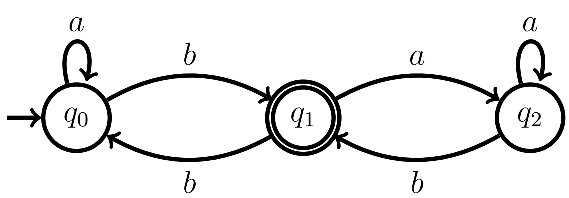
(c) Give an algorithm for the following operation:

INPUT: A state q of the master automaton.
OUTPUT: State r of the master automaton such that $L(r) = L(q)^R$.

(d) Give an algorithm for the following operation:

INPUT: A DFA A over alphabet Σ, and $k \in \mathbb{N}$.
OUTPUT: State q of the master automaton over Σ such that $L(q) = L(A) \cap \Sigma^k$.

Apply your algorithm on the following DFA with $k = 3$:



Exercise 8.3

Let $k \in \mathbb{N}_{>0}$. Let flip : $\{0,1\}^k \to \{0,1\}^k$ be the function that inverts the bits of its input, e.g. flip(010) = 101. Let val : $\{0,1\}^k \to \mathbb{N}$ be such that val(w) is the number represented by w in the *least significant bit first* encoding.

(a) Describe the minimal transducer that accepts

$$L_k = \left\{[x,y] \in (\{0,1\} \times \{0,1\})^k : \mathrm{val}(y) = \mathrm{val}(\mathrm{flip}(x)) + 1 \bmod 2^k\right\}.$$

(b) Build the state r of the master transducer for L_3, and the state q of the master automaton for $\{010, 110\}$.

(c) Adapt the algorithm *pre* seen in class to compute *post*(r, q).

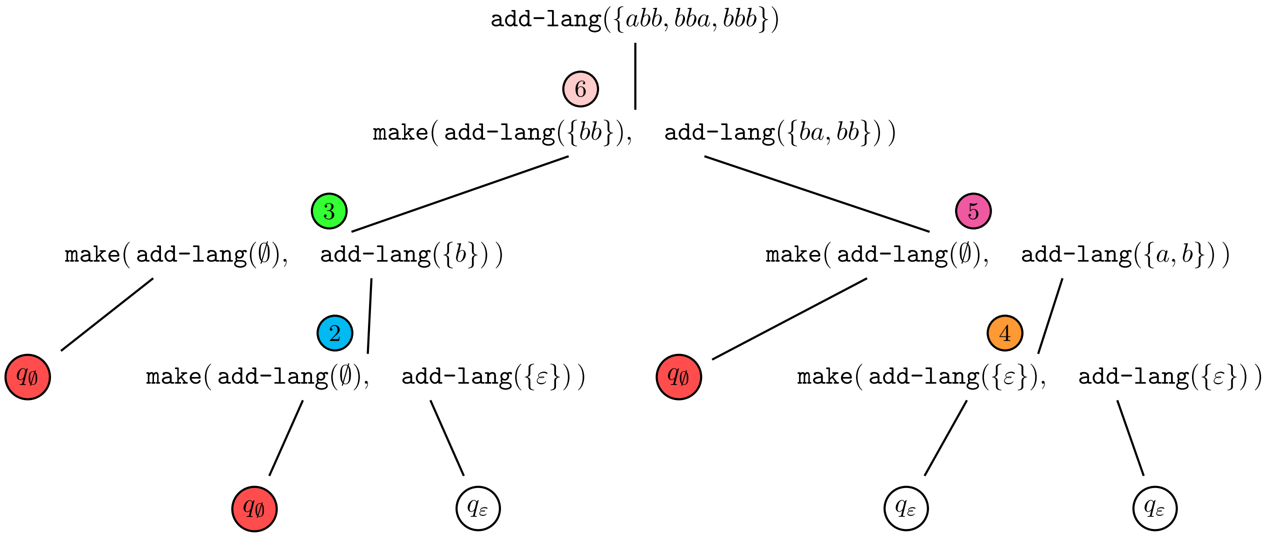
Solution 8.1

(a)

Input: A fixed-length language $L \subseteq \Sigma^k$ described explicitely by a set of words.
Output: State q of the master automaton over Σ such that $L(q) = L$.

add-lang(L) :
    if $L = \emptyset$ **then**
        return $q_\emptyset$
    else if $L = \{\varepsilon\}$ **then**
        return q_ε
    else
        for $a \in \Sigma$ **do**
            $L^a \leftarrow \{u : au \in L\}$
            $s_a \leftarrow$ add-lang(L^a)
        return make(s)

(b) Executing add-lang(L_1) yields the following computation tree:



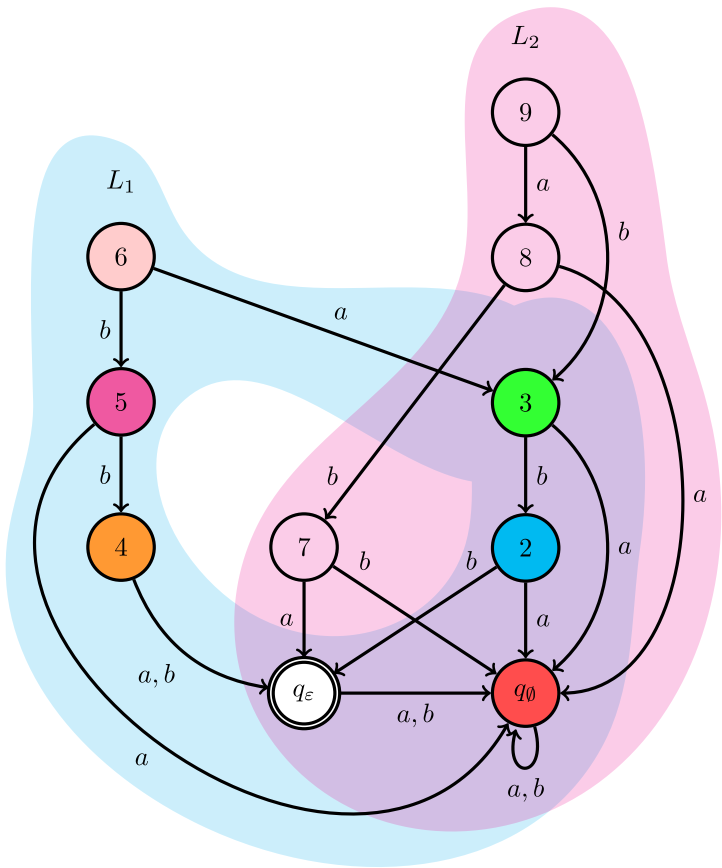
The table obtained after the execution is as follows:

Ident.	a-succ	b-succ
2	$q_\emptyset$	q_ε
3	$q_\emptyset$	2
4	q_ε	q_ε
5	$q_\emptyset$	4
6	3	5

Calling add-lang(L_2) adds the following rows to the table and returns 9:

Ident.	a-succ	b-succ
7	q_ε	$q_\emptyset$
8	$q_\emptyset$	7
9	8	3

The resulting master automaton fragment is:

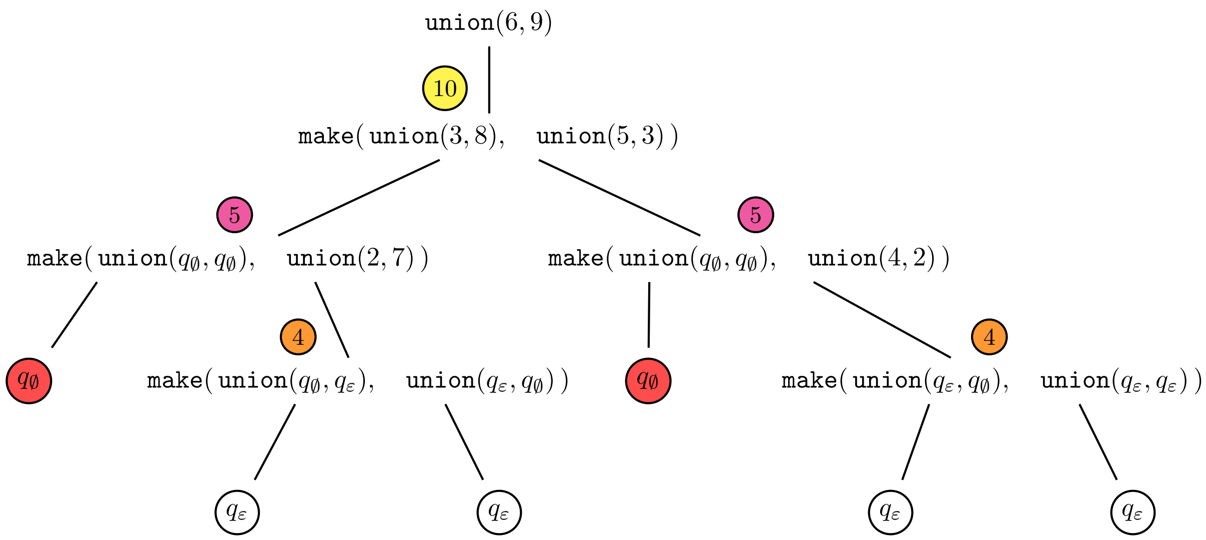


(c) Let us first adapt the algorithm for intersection to obtain an algorithm for union:

Input: States p and q of same length of the master automaton.
Output: State r of the master automaton such that $L(r) = L(p) \cup L(q)$.

`union`(p, q) *:*
  if $G(p, q)$ is not empty **then**
    return $G(p, q)$
  else if $p = q_\emptyset$ and $q = q_\emptyset$ **then**
    return $q_\emptyset$
  else if $p = q_\varepsilon$ or $q = q_\varepsilon$ **then**
    return q_ε
  else
    for $a \in \Sigma$ **do**
      $s_a \leftarrow$ `union`(p^a, q^a)
    $G(p, q) \leftarrow$ `make`(s)
    return $G(p, q)$

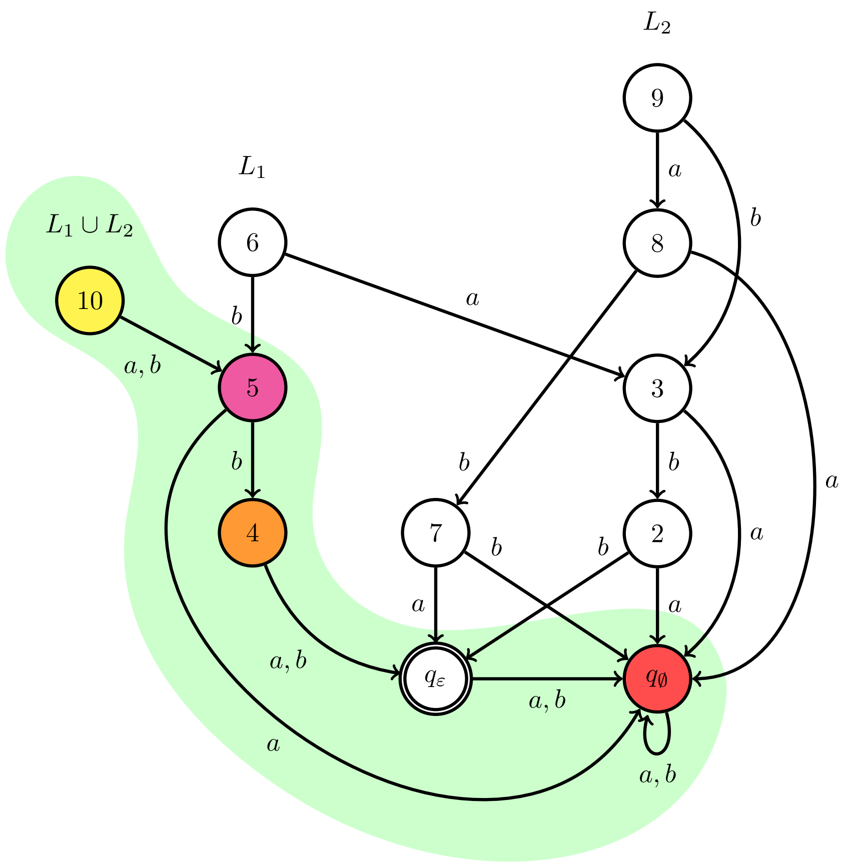
Executing `union`$(6, 9)$ yields the following computation tree:



Calling `union`$(6, 9)$ adds the following row to the table and returns 10:

Ident.	a-succ	b-succ
10	5	5

The new fragment of the master automaton is:



★ Note that `union` could be slightly improved by returning q whenever $p = q$, and by updating $G(q, p)$ at the same time as $G(p, q)$.

(d) The kernels are:

$$\langle L_1 \rangle = L_1,$$
$$\langle L_2 \rangle = L_2,$$
$$\langle L_1 \cup L_2 \rangle = \{ba, bb\}.$$

Solution 8.2

(a) Let L and L' be fixed-length languages. The following holds:

$$L \cdot L' = \begin{cases} \emptyset & \text{if } L = \emptyset, \\ L' & \text{if } L = \{\varepsilon\}, \\ \bigcup_{a \in \Sigma} a \cdot L^a \cdot L' & \text{otherwise.} \end{cases}$$

These identities give rise to the following algorithm:

Input: States p and q of the master automaton.
Output: State r of the master automaton such that $L(r) = L(p) \cdot L(q)$.

```
concat(p, q):
    if G(p, q) is not empty then
        return G(p, q)
    else if p = q_∅ then
        return q_∅
    else if p = q_ε then
        return q
    else
        for a ∈ Σ do
            s_a ← concat(p^a, q)
        G(p, q) ← make(s)
        return G(p, q)
```

(b) Let L be a fixed-length language and let h be a coding. The following holds:

$$h(L) = \begin{cases} \emptyset & \text{if } L = \emptyset, \\ \{\varepsilon\} & \text{if } L = \{\varepsilon\}, \\ \bigcup\limits_{a \in \Sigma} h(a) \cdot L^a & \text{otherwise.} \end{cases}$$

These identities give rise to the following algorithm:

Input: A state q of the master automaton and a coding h.
Output: State r of the master automaton such that $L(r) = \{h(w) : w \in L(q)\}$.

`coding`(q, h) *:*
  if $G(q)$ is not empty **then**
    return $G(q)$
  else if $q = q_\emptyset$ **then**
    return $q_\emptyset$
  else if $q = q_\varepsilon$ **then**
    return q_ε
  else
    $p \leftarrow q_\emptyset$
    for $a \in \Sigma$ **do**
      $r \leftarrow$ `coding`(q^a, h)
      $s_{h(a)} \leftarrow r$
      $s_b \leftarrow q_\emptyset$ for every $b \neq h(a)$
      $p \leftarrow$ `union`$(p,$ `make`$(s))$
    $G(q) \leftarrow p$
    return $G(q)$

The above algorithm makes use of `union` because the coding may be the same for distinct letters, i.e. $h(a) = h(b)$ for $a \neq b$ is possible. However, if the coding is a permutation, then this is not possible, and thus each letter maps to a unique residual. Therefore, the algorithm can be adapted as follows:

Input: A state q of the master automaton and a coding h which is a permutation.
Output: State r of the master automaton such that $L(r) = \{h(w) : w \in L(q)\}$.

`coding-permutation`(q, h) *:*
  if $G(q)$ is not empty **then**
    return $G(q)$
  else if $q = q_\emptyset$ **then**
    return $q_\emptyset$
  else if $q = q_\varepsilon$ **then**
    return q_ε
  else
    for $a \in \Sigma$ **do**
      $s_{h(a)} \leftarrow$ `coding-permutation`(q^a, h)
    $G(q) \leftarrow$ `make`(s)
    return $G(q)$

(c) Let L be a fixed-length language. The following holds:

$$L^R = \begin{cases} \emptyset & \text{if } L = \emptyset, \\ \{\varepsilon\} & \text{if } L = \{\varepsilon\}, \\ \bigcup_{a \in \Sigma} (L^a)^R \cdot a & \text{otherwise.} \end{cases}$$

These identities give rise to the following algorithm:

Input: A state q of the master automaton.
Output: State r of the master automaton such that $L(r) = L(q)^R$.

```
reverse(q):
    if G(q) is not empty then
        return G(q)
    else if q = q_∅ then
        return q_∅
    else if q = q_ε then
        return q_ε
    else
        p ← q_∅
        for a ∈ Σ do
            s_a ← q_ε
            s_b ← q_∅ for every b ≠ a
            r ← concat(reverse(q^a), make(s))
            p ← union(p, r)
        G(q) ← p
        return G(q)
```

(d) Let A be a DFA and let $k \in \mathbb{N}$. The following holds:

$$L(A) \cap \Sigma^k = \begin{cases} \emptyset & \text{if } k = 0 \text{ and } \varepsilon \notin L(A), \\ \{\varepsilon\} & \text{if } k = 0 \text{ and } \varepsilon \in L(A), \\ \bigcup_{a \in \Sigma} a \cdot (L(A)^a \cap \Sigma^{k-1}) & \text{otherwise.} \end{cases}$$

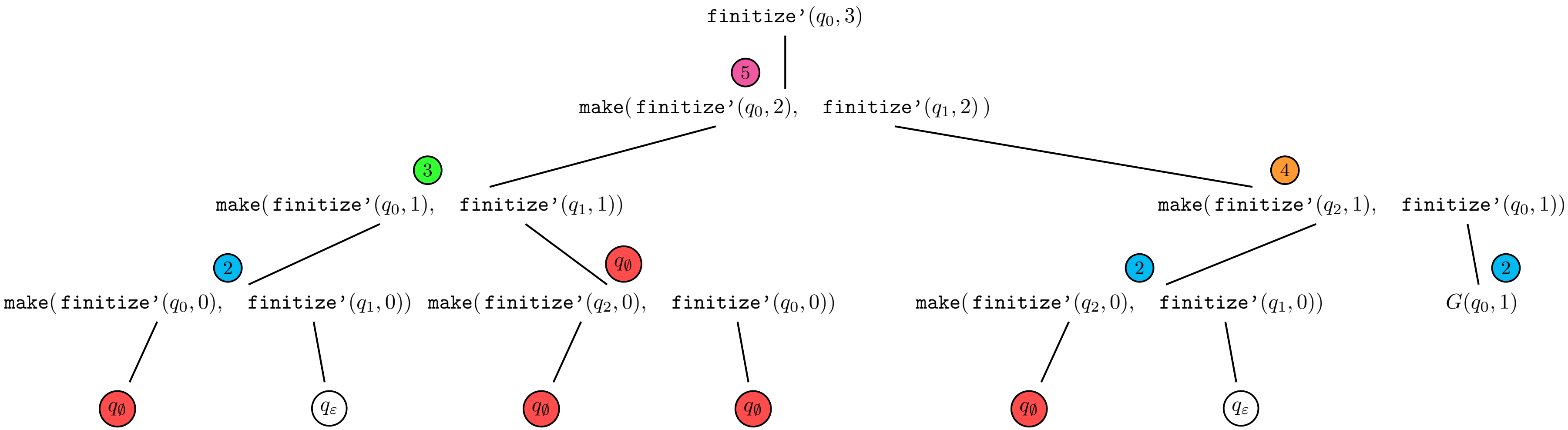
These identities give rise to the following algorithm:

Input: A DFA A over alphabet Σ, and $k \in \mathbb{N}$.
Output: State q of the master automaton over Σ such that $L(q) = L(A) \cap \Sigma^k$.

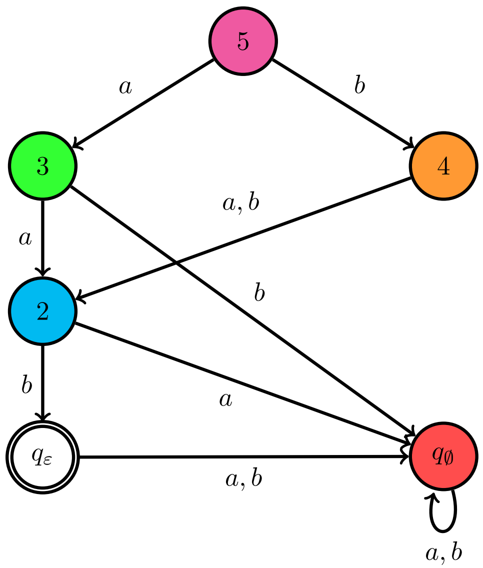
```
finitize(A, k):
    (Q, q_0, Σ, δ, F) ← A
    return finitize'(q_0, k)

finitize'(q, k):
    if G(q, k) is not empty then
        return G(q, k)
    else if k = 0 and q ∉ F then
        return q_∅
    else if k = 0 and q ∈ F then
        return q_ε
    else
        for a ∈ Σ do
            s_a ← finitize'(δ(q, a), k − 1)
        G(q, k) ← make(s)
        return G(q, k)
```

Executing `finitize`$(A, 3)$ calls `finitize'`$(q_0, 3)$ which yields the following computation tree:

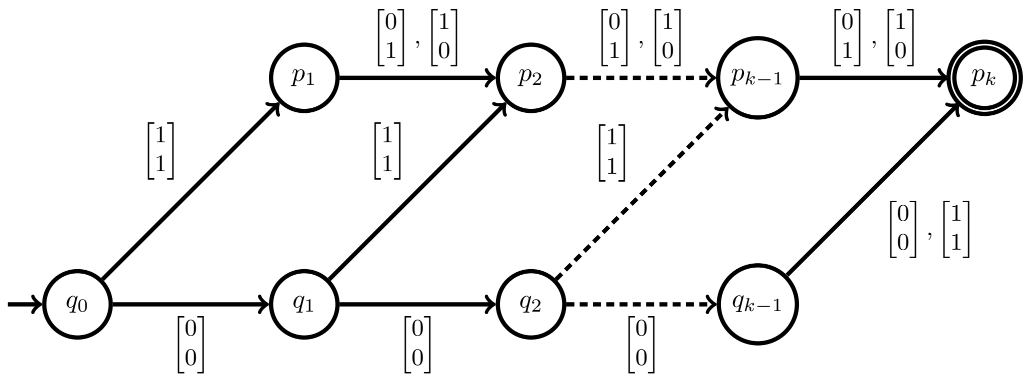


State 5 of the following master automaton fragment accepts $L(A) \cap \{a, b\}^3 = \{aab, bab, bbb\}$:

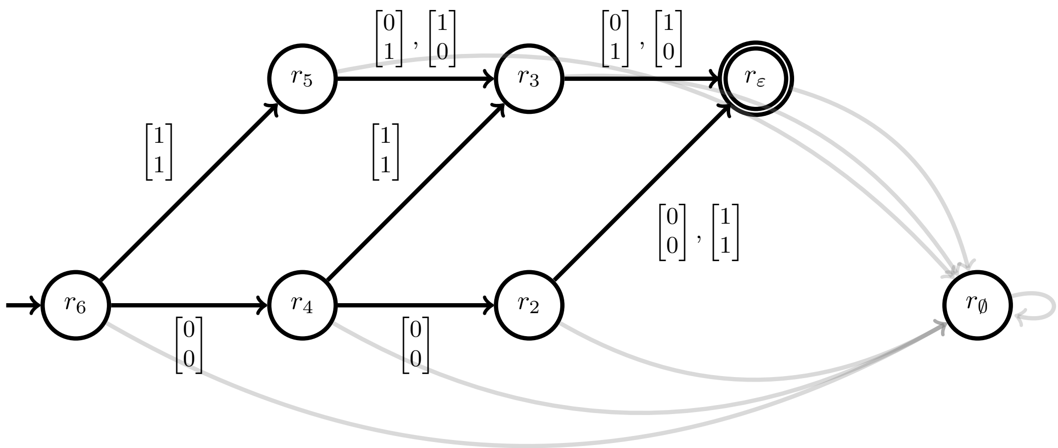


Solution 8.3

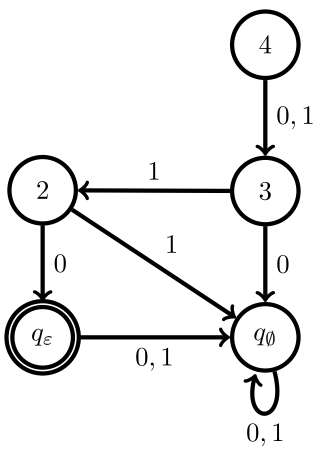
(a) Let $[x, y] \in L_k$. We may flip the bits of x at the same time as adding 1. If $x_1 = 1$, then $\neg x_1 = 0$, and hence adding 1 to val(flip(x)) results in $y_1 = 1$. Thus, for every $1 < i \leq k$, we have $y_i = \neg x_i$. If $x_1 = 0$, then $\neg x_1 = 1$. Adding 1 yields $y_1 = 0$ with a carry. This carry is propagated as long as $\neg x_i = 1$, and thus as long as $x_i = 0$. If some position j with $x_j = 1$ is encountered, the carry is "consumed", and we flip the remaining bits of x. These observations give rise to the following minimal transducer for L_k:



(b) The minimal transducer accepting L_3 is



State 4 of the following master automaton fragment accepts $\{010, 110\}$:



(c) We can establish the following identities similar to those obtained for *pre*:

$$
post_R(L) = \begin{cases} \emptyset & \text{if } R = \emptyset \text{ or } L = \emptyset, \\ \{\varepsilon\} & \text{if } R = \{[\varepsilon, \varepsilon]\} \text{ and } L = \{\varepsilon\}, \\ \displaystyle\bigcup_{a,b \in \Sigma} b \cdot post_{R^{[a,b]}}(L^a) & \text{otherwise.} \end{cases}
$$

To see that these identities hold, let $b \in \Sigma$ and $v \in \Sigma^k$ for some $k \in \mathbb{N}$. We have,

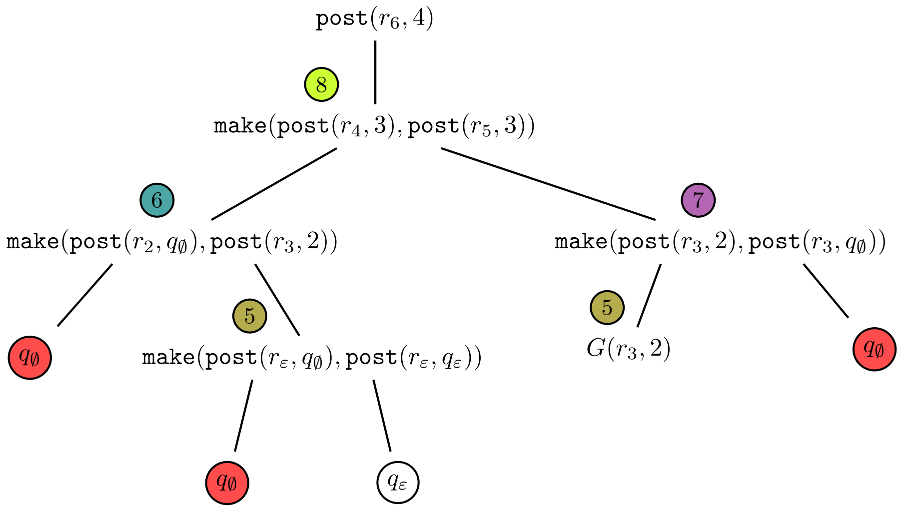
$$
\begin{aligned}
bv \in post_R(L) &\iff \exists a \in \Sigma, u \in \Sigma^k \text{ s.t. } au \in L \text{ and } [au, bv] \in R \\
&\iff \exists a \in \Sigma, u \in L^a \text{ s.t. } [au, bv] \in R \\
&\iff \exists a \in \Sigma, u \in L^a \text{ s.t. } [u, v] \in R^{[a,b]} \\
&\iff \exists a \in \Sigma \text{ s.t. } v \in Post_{R^{[a,b]}}(L^a) \\
&\iff v \in \bigcup_{a \in \Sigma} Post_{R^{[a,b]}}(L^a) \\
&\iff bv \in \bigcup_{a \in \Sigma} b \cdot Post_{R^{[a,b]}}(L^a).
\end{aligned}
$$

We obtain the following algorithm:

Input: A state r of the master transducer and a state q of the master automaton.
Output: State p of the master automaton such that $L(p) = Post_R(L)$ where $R = L(r)$ and $L = L(q)$.

`post`(r, q) *:*
    if $G(r, q)$ is not empty **then**
        return $G(r, q)$
    else if $r = r_\emptyset$ or $q = q_\emptyset$ **then**
        return $q_\emptyset$
    else if $r = r_\varepsilon$ and $q = q_\varepsilon$ **then**
        return q_ε
    else
        for $b \in \Sigma$ **do**
            $p \leftarrow q_\emptyset$
            for $a \in \Sigma$ **do**
                $p \leftarrow$ `union`$(p,$ `post`$(r^{[a,b]}, q^a))$
            $s_b \leftarrow p$
        $G(q, r) \leftarrow$ `make`(s)
        return $G(q, r)$

Note that the transducer for L_3 has some "strong" deterministic property. Indeed, for every state r and $b \in \{0,1\}$, if $r^{[a,b]} \neq r_\emptyset$ then $r^{[\neg a,b]} = r_\emptyset$. Hence, for a fixed $b \in \{0,1\}$, at most one term of the form "$\texttt{post}(r^{[a,b]}, q^a)$" can differ from $q_\emptyset$ at line 12 of the algorithm. Thus, unions made by the algorithm on this transducer are trivial, and executing $\texttt{post}(6,4)$ yields the following computation tree:



Calling $\texttt{post}(6,4)$ adds the following rows to the master automaton table and returns 8:

Ident.	0-succ	1-succ
5	$q_\emptyset$	q_ε
6	$q_\emptyset$	5
7	5	$q_\emptyset$
8	6	7

The resulting master automaton fragment: