## Automata and Formal Languages - Homework 7

Due 05.12.2017

## Exercise 7.1

Let val : $\{0,1\}^{*} \rightarrow \mathbb{N}$ be the function that associates to every word $w \in\{0,1\}^{*}$ the number $\operatorname{val}(w)$ represented by $w$ in the least significant bit first encoding.
(a) Give a transducer that doubles numbers, i.e. a transducer accepting

$$
L_{1}=\left\{[x, y] \in(\{0,1\} \times\{0,1\})^{*}: \operatorname{val}(y)=2 \cdot \operatorname{val}(x)\right\}
$$

(b) Give an algorithm that takes $k \in \mathbb{N}$ as input, and that produces a transducer $A_{k}$ accepting

$$
L_{k}=\left\{[x, y] \in(\{0,1\} \times\{0,1\})^{*}: \operatorname{val}(y)=2^{k} \cdot \operatorname{val}(x)\right\} .
$$

[Hint:
(c) Give a transducer for the addition of two numbers, i.e. a transducer accepting

$$
\left\{[x, y, z] \in(\{0,1\} \times\{0,1\} \times\{0,1\})^{*}: \operatorname{val}(z)=\operatorname{val}(x)+\operatorname{val}(y)\right\}
$$

(d) For every $k \in \mathbb{N}_{>0}$, let

$$
X_{k}=\left\{[x, y] \in(\{0,1\} \times\{0,1\})^{*}: \operatorname{val}(y)=k \cdot \operatorname{val}(x)\right\} .
$$

Sketch an algorithm that takes as input transducers $A$ and $B$, accepting respectively $X_{a}$ and $X_{b}$ for some $a, b \in \mathbb{N}_{>0}$, and that produces a transducer $C$ accepting $X_{a+b}$. [Hint:
(e) Let $k \in \mathbb{N}_{>0}$. Using (b) and (d), how can you build a transducer accepting $X_{k}$ ?
(f) Show that the following language has infinitely many residuals, and hence that it is not regular:

$$
\left\{[x, y] \in(\{0,1\} \times\{0,1\})^{*}: \operatorname{val}(y)=\operatorname{val}(x)^{2}\right\}
$$

## Exercise 7.2

Consider transducers whose transitions are labeled by elements of $(\Sigma \cup\{\varepsilon\}) \times\left(\Sigma^{*} \cup\{\varepsilon\}\right)$. Intuitively, each transition reads one or zero letter and writes a word of arbitrary length. Such a transducer can be used to perform operations on strings, e.g. upon reading singing in the rain it could write Singing In The Rain.

Sketch such $\varepsilon$-transducers for the following operations, each of which is informally defined by means of three examples. For each example, when the transducer reads the string on the left, it should write the string on the right. You may assume that the alphabet $\Sigma$ consists of $\{a, b, \ldots, z, A, B, \ldots, Z\}$, a whitespace symbol, and an end-of-line symbol. Moreover, you may assume that every string ends with an end-of-line symbol and contains no other occurrence of the end-of-line symbol.
(a)

| Input | Output |
| :--- | :---: |
| European Research Council | ERC |
| Technical University of Munich | TUM |
| FC Bayern | FCB |

(b)

| Input | Output |
| :--- | :--- |
| Finite automata rule | Finite automata rule |
| Transducers are fun | Transducers are fun |
| regular expressions are the best | regular expressions are the best |

(c)

| Input | Output |
| :--- | :--- |
| Hello world | HHEELLOO WWOORRLLDD |
| This is Sparta | TTHHIISS IISS SSPPAARRTTAA |
| over nine thousand | OOVVEERR NNIINNEE TTHHOOUUSSAANNDD |

(d) For this exercise, $\Sigma$ is extended with $\{, \cdot$,$\} .$

| Input | Output |
| :--- | :--- |
| Ada Lovelace | Lovelace, A. |
| Alan Turing | Turing, A. |
| Donald Knuth | Knuth, D. |

(e) For this exercise, $\Sigma$ is extended with $\{0,1, \ldots, 9,(),+$,$\} .$

| Input | Output |  |  |
| :--- | :--- | :--- | :--- |
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| $(00) 4989273452$ | +49 | 89 | 273452 |
| 273452 | +49 | 89 | 273452 |

