## Automata and Formal Languages - Homework 1

Due 21.10.2015

## Exercise 1.1

Go to http://www.jflap.org/ and download JFLAP. Run it and select the finite automata mode.
(a) Consider the language $C_{n}=\Sigma^{*} a \Sigma^{n-1}$ over $\Sigma=\{a, b\}$. Draw an NFA that recognizes $C_{3}$ and determinize it using JFLAP.
(b) Consider a similar language $D_{n}=\Sigma^{*} a(\Sigma \cup \varepsilon)^{n-1}$ over $\Sigma=\{a, b\}$. At least how many states does a DFA require to recognize $D_{n}$ ? Justify your answer.
(c) Let $L_{n}=\left\{a^{k} \mid k\right.$ is divisible by $n$ or $\left.n-1\right\}$ be a language over $\Sigma=\{a\}$. Draw an NFA $A$ that recognizes $L_{3}$.
(d) Use JFLAP to determinize $A$. How many states does $A$ have?
(e) Show that every DFA recognizing $L_{n}$ has at least $n(n-1)$ states.

## Exercise 1.2

Download a conversion game from https://www7.in.tum.de/tools/jflap-game/. Select the coversion game mode to play the game. Finish the following conversion types:
(a) Guess DFA from NFA, RE
(b) Guess NFA from RE
(c) Guess RE from DFA, NFA

## Exercise 1.3

Let $A$ be the following automaton:

(a) Transform the automaton $A$ into an equivalent regular expression, then transform this expression into an NFA (with $\varepsilon$-transitions), remove the $\varepsilon$-transitions, and determinize the automaton.
(b) Use JFLAP to perform the same transformations. Is there any difference?
(c) Use JFLAP to check that your resulting automaton is equivalent to the original one.

## Exercise 1.4

Given an alphabet $\Sigma$, we say that $w$ is a shuffle of words $u$ and $v$, if there exist $u_{i}, v_{i} \in \Sigma^{*}$ such that $u=u_{1} \cdots u_{k}$, $v=v_{1} \cdots v_{k}$, and $w=u_{1} v_{1} \cdots u_{k} v_{k}$.

Given languages $L_{1}$ and $L_{2}$, we define the shuffle of $L_{1}$ and $L_{2}$ as

$$
S\left(L_{1}, L_{2}\right)=\left\{w \mid \exists u \in L_{1}, v \in L_{2} \text { s.t. } w \text { is a shuffle of } u \text { and } v\right\}
$$

Show that if $L_{1}$ and $L_{2}$ are regular, then $S\left(L_{1}, L_{2}\right)$ is also regular.

