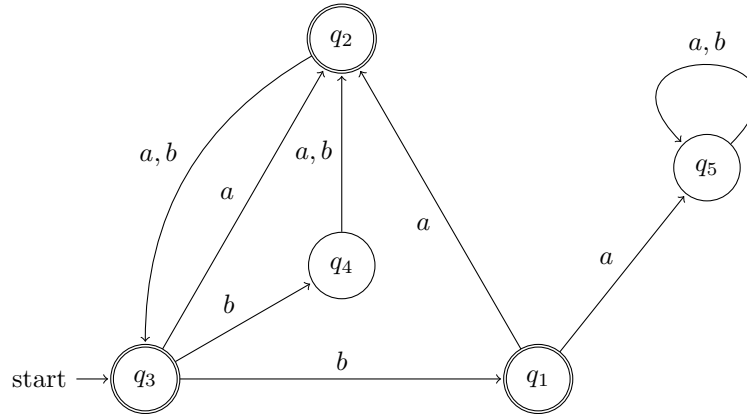


## Automata and Formal Languages – Homework 4

Due 10.11.2011.

### Exercise 4.1

Check whether the NFA depicted below recognizes  $\Sigma^*$  by means of the algorithm “UnivNFA” presented in the lecture.



### Exercise 4.2

We define the following languages over the alphabet  $\Sigma = \{a, b\}$ :

- $L_1$  is the set of all words where between any two occurrences of  $b$ 's there is at least one  $a$ .
- $L_2$  is the set of all words where every non-empty maximal sequence of consecutive  $a$ 's has odd length.
- $L_3$  is the set of all words where  $a$  occurs only at even positions.
- $L_4$  is the set of all words where  $a$  occurs only at odd positions.
- $L_5$  is the set of all words of odd length.
- $L_6$  is the set of all words with an even number of  $a$ 's.

*Remark:* For this exercise we assume that the first letter of a nonempty word is at position 1, e.g.,  $a \in L_4$ ,  $a \notin L_3$ .

Your task is to construct an FA, i.e., DFA or NFA or NFA- $\varepsilon$ , for

$$L := (L_1 \setminus L_2) \cup \overline{(L_3 \Delta L_4)} \cap L_5 \cap L_6 \text{ where } \Delta \text{ denotes the symmetric difference.}$$

while sticking to the following rules:

- You have to start from FAs for  $L_1, \dots, L_6$ .
- Any further automaton has to be constructed from already existing automata via an algorithm introduced in the lecture, e.g. Comp, BinOp, UnionNFA, NFAtoDFA, etc.

Try to find an order on the construction steps which yields an FA for  $L$  with as few states as possible.

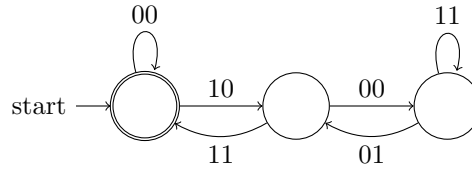
### Exercise 4.3

Find a family  $\{\mathcal{A}_n\}_{n=1}^\infty$  of NFAs with  $O(n)$  states such that *every* NFA recognizing the complement of  $\mathcal{L}(\mathcal{A}_n)$  has at least  $2^n$  states.

Hint: It is also possible to adapt the standard example  $L = \{ww \mid w \in \{a, b\}^*\}$ . (Here the complement of  $L$  is recognizable by push-down automata, whereas  $L$  is not.)

### Exercise 4.4

Consider the following FA  $\mathcal{A}$  over the alphabet  $\{00, 01, 10, 11\}$ :



W.r.t. the msbf encoding, we may interpret any word  $w \in \{00, 01, 10, 11\}^*$  as a pair of natural numbers  $(X(w), Y(w)) \in \mathbb{N}_0 \times \mathbb{N}_0$ . *Example:* (Underlined letters correspond to  $Y(w)$ .)

$$w = (00)^k 001011 \rightarrow (00)^k \underline{00} \underline{10} \underline{11} \rightarrow (0^k 011, 0^k 001) \rightarrow (3, 1) = (X(w), Y(w))$$

- (a) Show that  $w \in \mathcal{L}(\mathcal{A})$  iff  $X(w) = 3 \cdot Y(w)$ .
- (b) Construct the minimal DFA representing the language  $\{w \in \{0, 1\}^* \mid \text{msbf}^{-1}(w) \text{ is divisible by } 3\}$ .