Technische Universität München Theoretische Informatik Prof. J. Esparza / Dr.-Ing. J. Kreiker

# Exercises "Automata and Formal Languages"

## Exercise 6.1

Implement the crossing sequence construction and apply your implementation to the 2DFA of Exercise 5.3.

### Exercise 6.2

Define the n-ary asynchronous product of n NFAs, such that product states are n-element sets rather than n-tuples! Show that the thus defined n-ary product is associative and commutative.

### Exercise 6.3

In this exercise you are supposed to design a k-bit binary counter, C, of k communicating automata  $A_0, \ldots, A_{k-1}$ , such that  $C = A_0 || \ldots || A_{k-1}$ . Each automaton represents one bit. The alphabet of automaton  $A_i$  must contain (at least) the letters '0[i]' and '1[i]'. If the letter  $a_i[i]$  is read, it means that the *i*-th bit currently has value  $a_i$ . Whenever a letter '*inc*' is read, the number represented by the automata must be incremented by 1 (by a number of communications among the k automata). Design the automata  $A_i$  such that C satisfies the following specification:

- $\mathcal{L}(C) \mid_{\{inc\}} = inc^*$
- For all  $j \ge 0$ :  $inc^j a_0[0]a_1[1] \dots a_{k-1}[k-1] \in \mathcal{L}(C) \mid_{\{inc,0[i],1[i] \mid 0 \le i < k\}}$  if and only if  $j \mod 2^k = \sum_{i=0}^{k-1} a_i[i]2^i$

#### Exercise 6.4

In this exercise you are supposed to design a railway controller. Consider a circular railway of 8 segments:  $0 \to 1 \to \ldots \to 7 \to 0$ . There are three trains modeled by the automata  $T_0$ ,  $T_1$ , and  $T_2$ . Each automaton  $T_i$  has the states  $Q_i = \{q_{i,0}, \ldots, q_{i,7}\}$ , the alphabet  $\Sigma_i = \{enter[i, j] \mid 0 \le j \le 7\}$ , the transition relation  $\delta_i = \{(q_{i,j}, enter[i, j +_8 1], q_{i,j+s1}) \mid 0 \le j \le 7\}$ , and the initial state  $q_{i,2i}$ , where  $+_8$  is addition modulo 8.

Define a family  $C_1, \ldots C_k$  of controllers to make sure that two trains can never be on adjacent segments, while all trains can eventually move, if they want to. Formally, the product  $R = C_1 || \ldots ||C_k|| T_1 ||T_2|| T_3$  shall satisfy the following specification:

- For i = 0, 1, 2:  $\mathcal{L}(R) \mid_{\Sigma_i} = (enter[i, 2i]enter[i, 2i +_8 1] \dots enter[i, 2i +_8 7])^*$
- For all  $wenter[i, j]enter[i', j']w' \in \mathcal{L}(R) \mid_{\{enter[i, j] \mid i=0, 1, 2; 0 \le j \le 7\}}$  holds that  $i' \ne i$  implies  $|j j'| \notin \{0, 1, 7\}$

*Hint:* Typical values of k could be 1, 3, or 8 (where 8 is preferred).