

## Exercises “Automata and Formal Languages”

### Exercise 1.1

Let  $r$  be the regular expression  $((0 + 1)(0 + 1))^*$  over  $\Sigma = \{0, 1\}$ , where  $+$  denotes choice.

- Describe  $\mathcal{L}(r)$  in words.
- Give a regular expression  $r'$  such that  $\mathcal{L}(r') = \overline{\mathcal{L}(r)}$ , where  $\overline{L} = \Sigma^* \setminus L$ .
- *Construct* a regular expression  $r'$  such that  $\mathcal{L}(r') = \overline{\mathcal{L}(r)}$  using constructions learnt in class.

### Exercise 1.2

Let  $\Sigma$  be a finite alphabet. A function  $h : \Sigma \rightarrow \Sigma^*$  is called a *homomorphism*. We define:

$$\begin{aligned} h(\sigma_1 \dots \sigma_n) &= h(\sigma_1) \dots h(\sigma_n) && \text{for } \sigma_1, \dots, \sigma_n \in \Sigma \\ h(L) &= \{h(w) \mid w \in L\} && \text{for } L \subseteq \Sigma^* \\ h^{-1}(L) &= \{w \mid h(w) \in L\} && \text{for } L \subseteq \Sigma^* \end{aligned}$$

Let  $L \subseteq \Sigma^*$  be regular and let  $h : \Sigma \rightarrow \Sigma^*$  be a homomorphism.

- Show that  $h(L)$  regular.
- Show that  $h^{-1}(L)$  is regular.
- Use *only* closure properties and the fact that  $\{0^n 1^n \mid n \geq 0\}$  is not regular to show that the following languages over  $\{0, 1, 2\}$  are not regular:
  - $\{0^n 1^n 2^m \mid n, m \geq 0\}$
  - $\{0^{2n} 12^{3n} \mid n \geq 0\}$
  - $\{w \mid |w|_0 = |w|_2\}$ , where  $|w|_\sigma$  denotes the number of occurrences of letter  $\sigma$  in word  $w$ .

### Exercise 1.3

Since regular languages are closed under complement and intersection, we may extend regular expressions with such primitives: If  $r$  and  $r'$  are regular expressions, then also  $\bar{r}$  and  $r \cap r'$  are regular expressions. A language  $L \subseteq \Sigma^*$  is called *star-free*, iff there exists an extended regular expression  $r$  *without* a Kleene star such that  $L = \mathcal{L}(r)$ . For example,  $\Sigma^*$  is star-free, because it is the same as  $\mathcal{L}(\bar{\emptyset})$ . Show that  $\mathcal{L}((01 + 10)^*)$  is star-free.

### Exercise 1.4

An NFA  $(Q, \Sigma, \delta, q_0, F)$  (where  $\delta : Q \times \Sigma \rightarrow 2^Q$ ) is called *complete*, iff  $|\delta(q, a)| \geq 1$  for all  $q \in Q$  and  $a \in \Sigma$ . We define the *or-product* of two NFAs  $A_i = (Q_i, \Sigma, \delta_i, q_{0i}, F_i)$  ( $i = 1, 2$ ) to be  $A_1 \otimes A_2 = (Q_1 \times Q_2, \Sigma, \delta, (q_{01}, q_{02}), F)$ , where  $\delta$  is as in  $A_1 \times A_2$  and where  $F = \{(q_1, q_2) \mid q_1 \in F_1 \text{ or } q_2 \in F_2\}$ .

- Prove: If  $A_1$  and  $A_2$  are complete NFAs, then  $\mathcal{L}(A_1 \otimes A_2) = \mathcal{L}(A_1) \cup \mathcal{L}(A_2)$ .
- Completeness is thus sufficient for  $\mathcal{L}(A_1 \otimes A_2) = \mathcal{L}(A_1) \cup \mathcal{L}(A_2)$ . Is it also necessary?